## III. Recurrent Neural Networks

## A.

## The Hopfield Network

## Typical Artificial Neuron



## Typical Artificial Neuron



## Equations

Net input:

$$
\begin{aligned}
h_{i} & =\left(\sum_{j=1}^{n} w_{i j} s_{j}\right)-\theta \\
\mathbf{h} & =\mathbf{W} \mathbf{s}-\theta
\end{aligned}
$$

New neural state:

$$
\begin{aligned}
s_{i}^{\prime} & =\sigma\left(h_{i}\right) \\
\mathbf{s}^{\prime} & =\sigma(\mathbf{h})
\end{aligned}
$$

## Hopfield Network

- Symmetric weights: $w_{i j}=w_{j i}$
- No self-action: $w_{i i}=0$
- Zero threshold: $\theta=0$
- Bipolar states: $s_{i} \in\{-1,+1\}$
- Discontinuous bipolar activation function:

$$
\sigma(h)=\operatorname{sgn}(h)= \begin{cases}-1, & h<0 \\ +1, & h>0\end{cases}
$$

## What to do about $h=0$ ?

- There are several options:
- $\sigma(0)=+1$
- $\sigma(0)=-1$
- $\sigma(0)=-1$ or +1 with equal probability
- $h_{i}=0 \Rightarrow$ no state change $\left(s_{i}^{\prime}=s_{i}\right)$
- Not much difference, but be consistent
- Last option is slightly preferable, since symmetric


## Positive Coupling

- Positive sense (sign)
- Large strength



## Negative Coupling

- Negative sense (sign)
- Large strength



## Weak Coupling

- Either sense (sign)
- Little strength



## State $=-1 \&$ Local Field $<0$



## State $=-1 \&$ Local Field $>0$



## State Reverses



## State $=+1 \&$ Local Field $>0$



## State $=+1 \&$ Local Field $<0$



## State Reverses



# NetLogo Demonstration of Hopfield State Updating 

Run Hopfield-update.nlogo

## Hopfield Net as Soft Constraint Satisfaction System

- States of neurons as yes/no decisions
- Weights represent soft constraints between decisions
- hard constraints must be respected
- soft constraints have degrees of importance
- Decisions change to better respect constraints
- Is there an optimal set of decisions that best respects all constraints?


# Demonstration of Hopfield Net Dynamics I 

Run Hopfield-dynamics.nlogo

## Convergence

- Does such a system converge to a stable state?
- Under what conditions does it converge?
- There is a sense in which each step relaxes the "tension" in the system
- But could a relaxation of one neuron lead to greater tension in other places?


## Quantifying "Tension"

- If $w_{i j}>0$, then $s_{i}$ and $s_{j}$ want to have the same sign $\left(s_{i} s_{j}=+1\right)$
- If $w_{i j}<0$, then $s_{i}$ and $s_{j}$ want to have opposite signs $\left(s_{i} s_{j}=-1\right)$
- If $w_{i j}=0$, their signs are independent
- Strength of interaction varies with $\left|w_{i j}\right|$
- Define disharmony ("tension") $D_{i j}$ between neurons $i$ and $j$ :
$D_{i j}=-s_{i} w_{i j} s_{j}$
$D_{i j}<0 \Rightarrow$ they are happy
$D_{i j}>0 \Rightarrow$ they are unhappy


## Total Energy of System

The "energy" of the system is the total "tension" (disharmony) in it:

$$
\begin{aligned}
E\{\mathbf{s}\} & =\sum_{\langle i j\rangle} D_{i j}=-\sum_{\langle i j\rangle} s_{i} w_{i j} s_{j} \\
& =-\frac{1}{2} \sum_{i} \sum_{j \neq i} s_{i} w_{i j} s_{j} \\
& =-\frac{1}{2} \sum_{i} \sum_{j} s_{i} w_{i j} s_{j} \\
& =-\frac{1}{2} \mathbf{s}^{\mathrm{T}} \mathbf{W} \mathbf{S}
\end{aligned}
$$

## Review of Some Vector Notation

$$
\begin{aligned}
& \mathbf{x}=\left(\begin{array}{c}
x_{1} \\
\vdots \\
x_{n}
\end{array}\right)=\left(\begin{array}{lll}
x_{1} & \cdots & x_{n}
\end{array}\right)^{\mathrm{T}} \\
& \mathbf{x}^{\mathrm{T}} \mathbf{y}=\sum_{i=1}^{n} x_{i} y_{i}=\mathbf{x} \cdot \mathbf{y}
\end{aligned}
$$

(column vectors)

$$
\mathbf{x y}^{\mathrm{T}}=\left(\begin{array}{ccc}
x_{1} y_{1} & \cdots & x_{1} y_{n} \\
\vdots & \ddots & \vdots \\
x_{m} y_{1} & \cdots & x_{m} y_{n}
\end{array}\right)
$$

$$
\begin{equation*}
\mathbf{x}^{\mathrm{T}} \mathbf{M y}=\sum_{i=1}^{m} \sum_{j=1}^{n} x_{i} M_{i j} y_{j} \tag{quadraticform}
\end{equation*}
$$

## Another View of Energy

The energy measures the disharmony of the neurons' states with their local fields (i.e. of opposite sign):

$$
\begin{aligned}
E\{\mathbf{s}\} & =-\frac{1}{2} \sum_{i} \sum_{j} s_{i} w_{i j} s_{j} \\
& =-\frac{1}{2} \sum_{i} s_{i} \sum_{j} w_{i j} s_{j} \\
& =-\frac{1}{2} \sum_{i} s_{i} h_{i} \\
& =-\frac{1}{2} \mathbf{s}^{\mathrm{T}} \mathbf{h}
\end{aligned}
$$

## Do State Changes Decrease Energy?

- Suppose that neuron $k$ changes state
- Change of energy:

$$
\begin{aligned}
\Delta E & =E\left\{\mathbf{s}^{\prime}\right\}-E\{\mathbf{s}\} \\
& =-\sum_{\langle i j\rangle} s_{i}^{\prime} w_{i j} s_{j}^{\prime}+\sum_{\langle i j\rangle} s_{i} w_{i j} s_{j} \\
& =-\sum_{j \neq k} s_{k}^{\prime} w_{k j} s_{j}+\sum_{j \neq k} s_{k} w_{k j} s_{j} \\
& =-\left(s_{k}^{\prime}-s_{k}\right) \sum_{j \neq k} w_{k j} s_{j} \\
& =-\Delta s_{k} h_{k} \\
& <0
\end{aligned}
$$

## Energy Does Not Increase

- In each step in which a neuron is considered for update:
$E\{\mathbf{s}(t+1)\}-E\{\mathbf{s}(t)\} \leq 0$
- Energy cannot increase
- Energy decreases if any neuron changes
- Must it stop?


## Proof of Convergence in Finite Time

- There is a minimum possible energy:
- The number of possible states $\mathbf{s} \in\{-1,+1\}^{n}$ is finite
- Hence $E_{\text {min }}=\min \left\{E(\mathbf{s}) \mid \mathbf{s} \in\{ \pm 1\}^{n}\right\}$ exists
- Must reach in a finite number of steps because only finite number of states


## Conclusion

- If we do asynchronous updating, the Hopfield net must reach a stable, minimum energy state in a finite number of updates
- This does not imply that it is a global minimum


## Lyapunov Functions

- A way of showing the convergence of discreteor continuous-time dynamical systems
- For discrete-time system:
- need a Lyapunov function $E$ ("energy" of the state)
$-E$ is bounded below $\left(E\{\mathbf{s}\}>E_{\text {min }}\right)$
- $\Delta E<(\Delta E)_{\max } \leq 0$ (energy decreases a certain minimum amount each step)
- then the system will converge in finite time
- Problem: finding a suitable Lyapunov function


## Example Limit Cycle with Synchronous Updating



## The Hopfield Energy Function is Even

- A function $f$ is odd if $f(-x)=-f(x)$, for all $x$
- A function $f$ is even if $f(-x)=f(x)$, for all $x$
- Observe:

$$
E\{-\mathbf{s}\}=-\frac{1}{2}(-\mathbf{s})^{\mathrm{T}} \mathbf{W}(-\mathbf{s})=-\frac{1}{2} \mathbf{s}^{\mathrm{T}} \mathbf{W} \mathbf{s}=E\{\mathbf{s}\}
$$

## Conceptual Picture of Descent on Energy Surface




## Energy Surface



## Energy

 Surface $+$Flow Lines


## Flow Lines

## Basins of

 Attraction

# Bipolar State 

 Space

# Demonstration of Hopfield Net Dynamics II 

## Run initialized Hopfield.nlogo

# Storing Memories as Attractors 



# Example of Pattern Restoration 



# Example of Pattern Restoration 



## Example of Pattern Restoration



## Example of Pattern Restoration



## Example of Pattern Restoration



## Example of Pattern Completion



## Example of Pattern Completion



## Example of Pattern Completion



## Example of Pattern Completion



## Example of Pattern Completion



## Example of Association

## Example of Association



## Example of Association

## Example of Association



## Example of Association



## Applications of Hopfield Memory

- Pattern restoration
- Pattern completion
- Pattern generalization
- Pattern association


## Hopfield Net for Optimization and for Associative Memory

- For optimization:
- we know the weights (couplings)
- we want to know the minima (solutions)
- For associative memory:
- we know the minima (retrieval states)
- we want to know the weights


## Hebb’s Rule

"When an axon of cell A is near enough to excite a cell B and repeatedly or persistently takes part in firing it, some growth or metabolic change takes place in one or both cells such that A's efficiency, as one of the cells firing B , is increased."
-Donald Hebb (The Organization of Behavior, 1949, p. 62)
"Neurons that fire together, wire together"

## Example of Hebbian Learning: Pattern Imprinted



## Example of Hebbian Learning: Partial Pattern Reconstruction



## Mathematical Model of Hebbian Learning for One Pattern

Let $W_{i j}=\left\{\begin{array}{cc}x_{i} x_{j}, & \text { if } i \neq j \\ 0, & \text { if } i=j\end{array}\right.$
Since $x_{i} x_{i}=x_{i}^{2}=1, \quad \mathbf{W}=\mathbf{x x}{ }^{\mathrm{T}}-\mathbf{I}$
For simplicity, we will include self-coupling:

$$
\mathbf{W}=\mathbf{x} \mathbf{x}^{\mathrm{T}}
$$

## A Single Imprinted Pattern is a Stable State

- Suppose $\mathbf{W}=\mathbf{x x}^{\mathrm{T}}$
- Then $\mathbf{h}=\mathbf{W} \mathbf{x}=\mathbf{x x}^{\mathrm{T}} \mathbf{x}=n \mathbf{x}$ since

$$
\mathbf{x}^{\mathrm{T}} \mathbf{x}=\sum_{i=1}^{n} x_{i}^{2}=\sum_{i=1}^{n}( \pm \mathbf{1})^{2}=n
$$

- Hence, if initial state is $\mathbf{s}=\mathbf{x}$, then new state is $\mathbf{s}^{\prime}=\operatorname{sgn}(n \mathbf{x})=\mathbf{x}$
- For this reason, scale $\mathbf{W}$ by $1 / n$
- May be other stable states (e.g.,- $-\mathbf{x}$ )


## Questions

- How big is the basin of attraction of the imprinted pattern?
- How many patterns can be imprinted?
- Are there unneeded spurious stable states?
- These issues will be addressed in the context of multiple imprinted patterns


## Imprinting Multiple Patterns

- Let $\mathbf{x}^{1}, \mathbf{x}^{2}, \ldots, \mathbf{x}^{p}$ be patterns to be imprinted
- Define the sum-of-outer-products matrix:

$$
\begin{aligned}
W_{i j} & =\frac{1}{n} \sum_{k=1}^{p} x_{i}^{k} x_{j}^{k} \\
\mathbf{W} & =\frac{1}{n} \sum_{k=1}^{p} \mathbf{x}^{k}\left(\mathbf{x}^{k}\right)^{\mathrm{T}}
\end{aligned}
$$

## Definition of Covariance

Consider samples $\left(x^{1}, y^{1}\right),\left(x^{2}, y^{2}\right), \ldots,\left(x^{N}, y^{N}\right)$
Let $\bar{x}=\left\langle x^{k}\right\rangle$ and $\bar{y}=\left\langle y^{k}\right\rangle$
Covariance of $x$ and $y$ values:

$$
\begin{aligned}
C_{x y} & =\left\langle\left(x^{k}-\bar{x}\right)\left(y^{k}-\bar{y}\right)\right\rangle \\
& =\left\langle x^{k} y^{k}-\bar{x} y^{k}-x^{k} \bar{y}+\bar{x} \cdot \bar{y}\right\rangle \\
& =\left\langle x^{k} y^{k}\right\rangle-\bar{x}\left\langle y^{k}\right\rangle-\left\langle x^{k}\right\rangle \bar{y}+\bar{x} \cdot \bar{y} \\
& =\left\langle x^{k} y^{k}\right\rangle-\bar{x} \cdot \bar{y}-\bar{x} \cdot \bar{y}+\bar{x} \cdot \bar{y} \\
C_{x y} & =\left\langle x^{k} y^{k}\right\rangle-\bar{x} \cdot \bar{y}
\end{aligned}
$$



## Weights \& the Covariance Matrix

Sample pattern vectors: $\mathbf{x}^{1}, \mathbf{x}^{2}, \ldots, \mathbf{x}^{p}$
Covariance of $i^{\text {th }}$ and $j^{\text {th }}$ components:

$$
C_{i j}=\left\langle x_{i}^{k} x_{j}^{k}\right\rangle-\overline{x_{i}} \cdot \overline{x_{j}}
$$

If $\forall i: \overline{x_{i}}=0 \quad( \pm 1$ equally likely in all positions) :

$$
\begin{aligned}
& C_{i j}=\left\langle x_{i}^{k} x_{j}^{k}\right\rangle=\frac{1}{p} \sum_{k=1}^{p} x_{i}^{k} x_{j}^{k} \\
& \therefore n \mathbf{W}=p \mathbf{C}
\end{aligned}
$$

## Characteristics of Hopfield Memory

- Distributed ("holographic")
- every pattern is stored in every location (weight)
- Robust
- correct retrieval in spite of noise or error in patterns
- correct operation in spite of considerable weight damage or noise


# Demonstration of Hopfield Net 

Run Malasri Hopfield Demo

## Stability of Imprinted Memories

- Suppose the state is one of the imprinted patterns $\mathbf{x}^{m}$
- Then: $\mathbf{h}=\mathbf{W} \mathbf{x}^{m}=\left[\frac{1}{n} \sum_{k} \mathbf{x}^{k}\left(\mathbf{x}^{k}\right)^{\mathrm{T}}\right] \mathbf{x}^{m}$

$$
\begin{aligned}
& =\frac{1}{n} \sum_{k} \mathbf{x}^{k}\left(\mathbf{x}^{k}\right)^{\mathrm{T}} \mathbf{x}^{m} \\
& =\frac{1}{n} \mathbf{x}^{m}\left(\mathbf{x}^{m}\right)^{\mathrm{T}} \mathbf{x}^{m}+\frac{1}{n} \sum_{k \neq m} \mathbf{x}^{k}\left(\mathbf{x}^{k}\right)^{\mathrm{T}} \mathbf{x}^{m} \\
& =\mathbf{x}^{m}+\frac{1}{n} \sum_{k \neq m}\left(\mathbf{x}^{k} \cdot \mathbf{x}^{m}\right) \mathbf{x}^{k}
\end{aligned}
$$

## Interpretation of Inner Products

- $\mathbf{x}^{k} \cdot \mathbf{x}^{m}=n$ if they are identical
- highly correlated
- $\mathbf{x}^{k} \cdot \mathbf{x}^{m}=-n$ if they are complementary - highly correlated (reversed)
- $\mathbf{x}^{k} \cdot \mathbf{x}^{m}=0$ if they are orthogonal - largely uncorrelated
- $\mathbf{x}^{k} \cdot \mathbf{x}^{m}$ measures the crosstalk between patterns $k$ and $m$


## Cosines and Inner products

$\mathbf{u} \cdot \mathbf{v}=\|\mathbf{u}\|\|\mathbf{v}\| \cos \theta_{\mathbf{u v}}$


If $\mathbf{u}$ is bipolar, then $\|\mathbf{u}\|^{2}=\mathbf{u} \cdot \mathbf{u}=n$
Hence, $\mathbf{u} \cdot \mathbf{v}=\sqrt{n} \sqrt{n} \cos \theta_{\mathbf{u v}}=n \cos \theta_{\mathbf{u v}}$

$$
\text { Hence } \mathbf{h}=\mathbf{x}^{m}+\sum_{k \neq m} \mathbf{x}^{k} \cos \theta_{k m}
$$

## Conditions for Stability

Stability of entire pattern:

$$
\mathbf{x}^{m}=\operatorname{sgn}\left(\mathbf{x}^{m}+\sum_{k \neq m} \mathbf{x}^{k} \cos \theta_{k m}\right)
$$

Stability of a single bit:

$$
x_{i}^{m}=\operatorname{sgn}\left(x_{i}^{m}+\sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}\right)
$$

## Sufficient Conditions for Instability (Case 1)

Suppose $x_{i}^{m}=-1$. Then unstable if :

$$
\begin{array}{r}
(-1)+\sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}>0 \\
\sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}>1
\end{array}
$$

## Sufficient Conditions for Instability (Case 2)

Suppose $x_{i}^{m}=+1$. Then unstable if :

$$
\begin{aligned}
(+1)+ & \sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}<0 \\
& \sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}<-1
\end{aligned}
$$

## Sufficient Conditions for Stability

$$
\left|\sum_{k \neq m} x_{i}^{k} \cos \theta_{k m}\right| \leq 1
$$

The crosstalk with the sought pattern must be sufficiently small

## Capacity of Hopfield Memory

- Depends on the patterns imprinted
- If orthogonal, $p_{\max }=n$
- but every state is stable $\Rightarrow$ trivial basins
- So $p_{\max }<n$
- Let load parameter $\alpha=p / n$


## Single Bit Stability Analysis

- For simplicity, suppose $\mathbf{x}^{k}$ are random
- Then $\mathbf{x}^{k} \cdot \mathbf{x}^{m}$ are sums of $n$ random $\pm 1$
- binomial distribution $\approx$ Gaussian
- in range $-n, \ldots,+n$
- with mean $\mu=0$
- and variance $\sigma^{2}=n$

- Probability sum $>t$ :

$$
\frac{1}{2}\left[1-\operatorname{erf}\left(\frac{t}{\sqrt{2 n}}\right)\right]
$$

[See "Review of Gaussian (Normal) Distributions" on course website]

## Approximation of Probability

Let crosstalk $C_{i}^{m}=\frac{1}{n} \sum_{k \neq m} x_{i}^{k}\left(\mathbf{x}^{k} \cdot \mathbf{x}^{m}\right)$
We want $\operatorname{Pr}\left\{C_{i}^{m}>1\right\}=\operatorname{Pr}\left\{n C_{i}^{m}>n\right\}$
Note : $n C_{i}^{m}=\sum_{\substack{k=1 \\ k \neq m}}^{p} \sum_{j=1}^{n} x_{i}^{k} x_{j}^{k} x_{j}^{m}$
A sum of $n(p-1) \approx n p$ random $\pm 1 \mathrm{~s}$
Variance $\sigma^{2}=n p$

## Probability of Bit Instability



## Tabulated Probability of Single-Bit Instability

| $P_{\text {error }}$ | $\alpha$ |
| :--- | :--- |
| $0.1 \%$ | 0.105 |
| $0.36 \%$ | 0.138 |
| $1 \%$ | 0.185 |
| $5 \%$ | 0.37 |
| $10 \%$ | 0.61 |

## Orthogonality of Random Bipolar Vectors of High Dimension

- $99.99 \%$ probability of being within $4 \sigma$ of mean
- It is $99.99 \%$ probable that random $n$-dimensional vectors will be within $\varepsilon=4 / \sqrt{ } n$ orthogonal
- Probability of being less orthogonal than $\varepsilon$ decreases exponentially with $n$
- The brain gets approximate orthogonality by assigning random high-dimensional vectors

$$
\operatorname{Pr}\{|\cos \theta|>\varepsilon\}=\operatorname{erfc}\left(\frac{\varepsilon \sqrt{n}}{\sqrt{2}}\right)
$$

$$
\approx \frac{1}{6} \exp \left(-\varepsilon^{2} n / 2\right)+\frac{1}{2} \exp \left(-2 \varepsilon^{2} n / 3\right)
$$

## Spurious Attractors

- Mixture states:
- sums or differences of odd numbers of retrieval states
- number increases combinatorially with $p$
- shallower, smaller basins
- basins of mixtures swamp basins of retrieval states $\Rightarrow$ overload
- useful as combinatorial generalizations?
- self-coupling generates spurious attractors
- Spin-glass states:
- not correlated with any finite number of imprinted patterns
- occur beyond overload because weights effectively random


## Basins of Mixture States



## Fraction of Unstable Imprints ( $n=100$ )



## Number of Stable Imprints ( $n=100$ )



## Number of Imprints with Basins of Indicated Size $(n=100)$



## Summary of Capacity Results

- Absolute limit: $p_{\max }<\alpha_{c} n=0.138 n$
- If a small number of errors in each pattern permitted: $p_{\text {max }} \propto n$
- If all or most patterns must be recalled perfectly: $p_{\text {max }} \propto n / \log n$
- Recall: all this analysis is based on random patterns
- Unrealistic, but sometimes can be arranged

