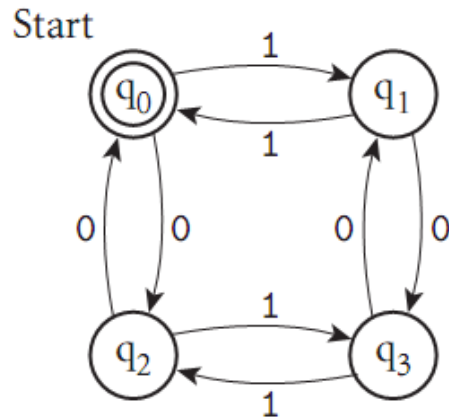


# Functional Programming Demonstration

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# A Review/Overview of Scheme

## Example program - Simulation of DFA



```
(define zero-one-even-dfa
  '(q0                                     ; start state
    (((q0 0) q2) ((q0 1) q1) ((q1 0) q3) ((q1 1) q0) ; transition fn
      ((q2 0) q0) ((q2 1) q3) ((q3 0) q1) ((q3 1) q2))
    (q0))) ; final states
```

**Figure 10.2** DFA to accept all strings of zeros and ones containing an even number of each. At the bottom of the figure is a representation of the machine as a Scheme data structure, using the conventions of Figure 10.1.

# A Review/Overview of Scheme

## Example program - Simulation of DFA

```
(define simulate
  (lambda (dfa input)
    (cons (current-state dfa)          ; start state
          (if (null? input)
              (if (infinal? dfa) '(accept) '(reject))
              (simulate (move dfa (car input)) (cdr input))))))

;; access functions for machine description:
(define current-state car)
(define transition-function cadr)
(define final-states caddr)
(define infinal?
  (lambda (dfa)
    (memq (current-state dfa) (final-states dfa))))

(define move
  (lambda (dfa symbol)
    (let ((cs (current-state dfa)) (trans (transition-function dfa)))
      (list
        (if (eq? cs 'error)
            'error
            (let ((pair (assoc (list cs symbol) trans)))
              (if pair (cadr pair) 'error))) ; new start state
        trans ; same transition function
        (final-states dfa)))) ; same final states
```

**Figure 10.1** Scheme program to simulate the actions of a DFA. Given a machine description and an input symbol  $i$ , function `move` searches for a transition labeled  $i$  from the start state to some new state  $s$ . It then returns a new machine with the same transition function and final states, but with  $s$  as its "start" state. The main function, `simulate`, tests to see if it is in a final state. If not, it passes the current machine description and the first symbol of input to `move`, and then calls itself recursively on the new machine and the remainder of the input. The functions `cadr` and `caddr` are defined as `(lambda (x) (car (cdr x)))` and `(lambda (x) (car (cdr (cdr x))))`, respectively. Scheme provides a large collection of such abbreviations.