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## Did anything change in the meantime?

- Parallel algorithm design and scheduling were already difficult tasks with homogeneous machines
- On heterogeneous platforms, it gets worse
- Patrick Geoffray went from kindergarten to Myricom but he's still a kid!
- He says that only embarrassingly parallel applications can be deployed on the grid
- Clearly, he is over optimistic! (3)


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## A nice little embarrassingly parallel application



- One (divisible load) application running on each cluster $\Rightarrow$ Which fraction of the job to delegate to other clusters?
- Different communication-to-computation ratios $\Rightarrow$ How to ensure fair scheduling and good resource utilization?


# Revisiting matrix product on heterogeneous platforms 

Jack Dongarra, Zhiao Shi, UT Knwoville
Jean-François Pineau, Yves Robert, Frédéric Vivien, ENS Lyon

Revisiting matr product on heteroge cos platforms

Eh wait!
Experiments are not ready?!

Jear ançois Pineau, Yves Robert, Frédéric Vivien, E Lyon

# Scheduling and Data Redistribution Strategies on Star Platforms 

Loris Marchal, Veronika Rehn,<br>Yves Robert and Frédéric Vivien

GRAAL team, LIP<br>École Normale Supérieure de Lyon

September 2006

## Outline

(1) Target problem

- Fully homogeneous platforms
- Bus platforms
- General platforms
(2) Simulations
(3) Divisible Loads Using the Multiport Switch-Model

4 Conclusion

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## Example

## Architecture




## Example

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## Example

## Architecture <br> 



## Example

## Architecture




## Framework

- Master-slave platforms
- New: Distributed loads


## Problem <br> Redistribution of data Goal: Minimize overall processing time

 Data models- Independent tasks - Divisible loads


## Framework

- Master-slave platforms
- New: Distributed loads


## Problem

Redistribution of data
Goal: Minimize overall processing time

- Independent tasks
- Divisible loads
- Master-slave platforms
- New: Distributed loads


## Problem

Redistribution of data
Goal: Minimize overall processing time
Data models

- Independent tasks
- Divisible loads


## Related Work

Independent tasks

- Application: BOINC (e.g. Einstein@home)
- NP-completeness for different task sizes

Divisible load theory

- Perfect parallel jobs
- Ontimal algorithms for video processing (Altilar, Paker)

Redistribution algorithms

- NP-completeness (Kremer)
- Optimality for particular cases: homogeneous ring topologies


## Related Work

## Independent tasks



Our approach


- Application: BOINC (e.g. Einstein@home)
- NP-completeness for different task sizes

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## Model

- Star network $S=P_{0}, P_{1}, \ldots, P_{m}$
- Communication cost $c_{i}$
- Computing power $w_{i}$
- Initial data $L_{i}$
- Independent and identical tasks
- Linear cost model
- Bidirectional one-port model
- Objective function

Minimize makespan

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## Best-Balance Algorithm - BBA

- Homogeneous communication links
- Homogeneous workers


Principle: Local optimization of current makespan

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## BBA- Optimality

## Theorem

Best-Balance Algorithm calculates an optimal schedule $S$ on a fully homogeneous star network.

## Moore-Based Binary-Search Algorithm - MBBSA

- Homogeneous communication links
- Heterogeneous workers
- Makespan M

Principle:

- Moore's algorithm
- Schedule within M
- Binary search

Moore's algorithm
Order the jobs by
. end if
end for
non-decreasing deadlines:


## Moore-Based Binary-Search Algorithm - MBBSA

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## Moore's algorithm

1: Order the jobs by non-decreasing deadlines:
$d_{1} \leq d_{2} \leq \cdots \leq d_{d}$
2: $\sigma \leftarrow \emptyset ; t \leftarrow 0$
3: for $i:=1$ to $n$ do
4: $\quad \sigma \leftarrow \sigma \cup\{i\}$
5: $\quad t \leftarrow t+w_{i}$
6: if $t>d_{i}$ then
7: $\quad$ Find job $j$ in $\sigma$ with largest $w_{j}$ value
8: $\quad \sigma \leftarrow \sigma \backslash\{j\}$
$t \leftarrow t-w_{j}$
10: end if
11: end for

## Bus platforms

## MBBSA- Phase $1+2$

Determination of senders and receivers


## Bus platforms

## MBBSA- Phase $1+2$

Determination of senders and receivers


## MBBSA- Phase 3

## Computation of deadlines



## MBBSA- Phase 4

Scheduling step


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Scheduling step


## MBBSA- Phase 4

Scheduling step


## MBBSA- Optimality

## Theorem

(i) MBBSA succeeds to build a schedule $\sigma$ for a given makespan $M$, if and only if there exists one.
(ii) Binary search algorithm returns in polynomial time an optimal schedule $\sigma$ for bus platforms (homogeneous communication links and heterogeneous workers).

## Dealing with fully heterogeneous platforms

Difficulty: Who is sender, who is receiver?
$\mathrm{M}=12$

| Worker | c | w | load |
| :--- | :---: | :---: | :---: |
| $P_{1}$ | 1 | 1 | 13 |
| $P_{2}$ | 8 | 1 | 13 |
| $P_{3}$ | 1 | 9 | 0 |
| $P_{4}$ | 1 | 10 | 0 |



## NP-completeness

Scheduling Problem for Master-Slave Tasks on a Star of Heterogeneous Processors

## Definition (SPMSTSHP)

Let $N$ be a star-network. Let $T$ be a deadline.
"Is it possible to redistribute tasks and process them in time $T$ ?".

## Theorem

$N P$-complete in the strong sense.

## Proof: Reduction to 3-partition

Proof: Reduction to 3-partition problem


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## Impact of Heterogeneity

| Platform type |  |  |
| :---: | :---: | :---: |
| Comm. | Comp. | Difficulty |
| Hom. | Hom. | simple greedy algorithm |
| Hom. | Het. | complicated algorithm |
| Het. | Hom. | ? |
| Het. | Het. | NP-strong |

## Heuristics

- BBA
- MBBSA
- R-BSA: Reversed-Binary Search Algorithm Combination of greedy algorithm and binary search


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## Simulations

## SimGRID:

## Simulator for distributed applications

- 4 platform types
- 1000 instances
- 10 workers
- Random variables
- $c_{i}: 1 . .100$
- $w_{i}: 1 . .100$
- $L_{i}: 0 . .50$


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## Trace Tests

## BBA



## MBBSA



## Distance from the Best Heuristic

Heterogeneous platform


## Distance from the Best Heuristic

Heterogeneous platform


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## Distance from the Best Heuristic

Heterogeneous platform


## Standard Deviation

| Platform type |  | Standard deviation |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Comm. | Comp. | BBA | MBBSA | R-BSA |
| Hom | Hom | 0 | 0 | 0.0107 |
| Hom | Het | 0.0006 | 0 | 0.0181 |
| Het | Hom | 0.4007 | 0.0208 | 0.0173 |
| Het | Het | 0.3516 | 0.0327 | 0.0284 |

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## Framework

## Star network



- Switch as master
- $m$ workers
- Computation speed $s_{i}$
- Bandwidth $b_{i}$
- Divisible load $\alpha_{i}$
- Linear cost model
- Overlapped unbounded switch model


## Redistribution Strategy

Goal: Every worker finishes at the same time


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## Solution for Divisible Loads

Imbalance of a worker $\delta_{i}$
Linear program

Minimize $T$,
UNDER THE CONSTRAINTS

## Fraction of load $f_{i, j}$

$$
\left\{\begin{array}{l}
(1 \mathrm{a}) \quad\left|\delta_{i}\right| \leq T \times b_{i} \\
(1 \mathrm{~b}) \\
\delta_{i} \geq \alpha_{i}-T \times s_{i} \\
(1 \mathrm{c}) \quad \sum_{i} \delta_{i}=0
\end{array}\right.
$$

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\end{array} \sum_{i} \delta_{i}=0\right.
$$

(1)

Fraction of load $f_{i, j}$

$$
f_{i, j}=\delta_{i} \times \frac{\delta_{j}}{\sum_{k \in R} \delta_{k}}=\delta_{i} \times \frac{\delta_{j}}{-L}
$$

Communication rate $\lambda_{i, j}$

$$
\lambda_{i, j}=\frac{f_{i, j}}{T_{0}}
$$

Computation rate $\gamma_{i, j}$

$$
\gamma_{i, j}=\frac{f_{i, j}}{T_{0}}
$$

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## Conclusion

Complete study of a difficult load-balancing problem
Scheduling and redistributing data on master-slave platforms Independent tasks:

- General case: Proof of NP-completeness in the strong sense
- Special platforms: Optimal algorithms
- Simulations: Verification of theoretical results

Divisible loads:

- Solution for general case: LP + analytical formulas


## Perspectives

Beyond the NP-completeness: Search for approximation algorithms Extension to dynamic master-slave platforms Extension to more general interconnection networks

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Extension to more general interconnection networks

## Recent work at our place

On-line scheduling heuristics for master-slave platforms

- Competitive ratios and inapproximability results
- Communication-aware heuristics

Collective communications

- Broadcast, multicast on heterogeneous clusters
- Resource selection for future MPI2 routines

Load-balancing

- Optimize BOINC-like applications
- Data redistribution strategies

Steady-state scheduling

- Multiple applications competing for resources
- Centralized vs fully distributed heuristics


## Scheduling for large-scale platforms

Assess the impact of new architectural characteristics

- Heterogeneity
- Irregular network topologies
- Hierarchy
- Variability (volability)

Inject static knowledge in a (mostly) dynamic environment

- Divisible loads vs bag of tasks
- Steady-state scheduling
- Resource selection

Evaluation

- Evaluate strategies through simulation
- SimGrid software co-developed with UCSD
- Large-scale experiments with Grid'5000

