

Lecture 13

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Project 2

- See course website under “Projects / Assignments”
- Due Oct. 25
- This project involves programming as well as some lengthy computations
- Start early!

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Digression: Time-Reversibility and the Physical Limits of Computation

Work done by:

- Rolf Landauer (1961)
- Charles Bennett (1973)
- Richard Feynman (1981–3)

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Entropy of Physical Systems

- Recall information content of N equally likely messages: $I_2 = \lg N$ bits.
- Can also use natural logs: $I_e = \ln N$ nats = $I_2 \ln 2$.
- To specify position & momentum of each particle of a physical system:
 $S = k \ln N = k I_2 \ln 2$.
 - k is Boltzmann’s constant
- This is the entropy of the system
 - entropies of 10 bits/atom are typical

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Thermodynamics of Recording One Bit



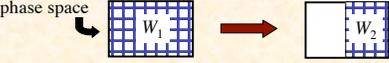
$$\Delta S = k \ln \frac{V_2}{V_1}$$

if $V_2 = V_1/2$, then $\Delta S = -k \ln 2$
 also, $\Delta F = kT \ln 2$

- ΔS derived by gas laws & classical thermodynamics
- Boltzmann constant: $k = 1.381 \times 10^{-23} \text{ J K}^{-1}$

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Entropy Change in Terms of Phase Space



- Let W = number of microstates corresponding to a macrostate
- Entropy $S = k \ln W$
- Then $\Delta S = k \ln W_2 - k \ln W_1 = k \ln (W_2 / W_1)$
- If $W_2 = W_1 / 2$, then $\Delta S = -k \ln 2$

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Information and Energy

random register energy initialized register
 fuel value = 0 energy fuel value = $nkT \ln 2$
 information = n bits information = 0 bits

- initialization equivalent to storing energy
- information and energy are complementary

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Minimum Energy for Irreversible Computation

2 bits 0 1 0 1 1 bit
 0 1 AND 0 1

- Loss of one bit of information
 - an irreversible operation (many to one)
- $\Delta S = -k \ln 2$
 - entropy decrease must be compensated by heat dissipation
- Minimum energy required: $\Delta F = kT \ln 2$
 - transistors: $\sim 10^8 kT$; RNA polymerase: $\sim 100 kT$

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Reversible Gates

- Can make dissipation arbitrarily small by using reversible gates
- All outputs must go somewhere
- Cannot *ever* throw information away
- The Fredkin CCN gate (“Controlled Controlled Not”) is reversible
 - can be used for constructing other gates

control lines $\left\{ \begin{array}{l} A \rightarrow A' = A \\ B \rightarrow B' = B \\ C \rightarrow C' = \begin{cases} \neg C & \text{if } A \wedge B \\ C & \text{otherwise} \end{cases} \end{array} \right.$

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Reversible Computer

input (m bits) M (all reversible gates) copy of input (m bits)
 zeroes (n bits) output (n bits)

- Reversible because get input back
- Only loss is resetting machine for next job
 - energy is proportional to n , number of output bits

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Summary: Energy Required for Reversible Computing

- There is no lower limit on the energy required for basic operations (gates, bit copying, etc.) provided:
 - it is done reversibly
 - it is done sufficiently slowly
- What is the fundamental relation between speed and energy dissipation?

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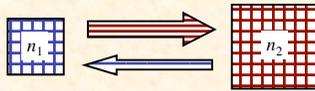
Energy and the Speed of Computation

ΔE E₁ E₂

- Let r be ratio of forward to backward rate
- Statistical mechanics shows: $kT \ln r = \Delta E$
- Greater “driving energy” \Rightarrow greater rate

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Entropy and the Speed of Computation



- Consider number of accessible microstates, n_1 and n_2
- Can show: $r = n_2 / n_1$
- Hence, $kT \ln r = kT (\ln n_2 - \ln n_1)$
 $= (S_2 - S_1)T = T\Delta S$

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Conclusions

- Entropy increase and energy dissipation can be made arbitrarily small by doing reversible computation
- However, the speed of computation is an exponential function of the driving energy or entropy increase

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Additional Information

1. Feynman, R. P. *Feynman Lectures on Computation*, ed. by A.J.G Hey & R.W. Allen. Perseus, 1996.
2. Hey, A.J.G. (ed.) *Feynman and Computation: Exploring the Limits of Computation*. Perseus, 1999.

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Real Ants

(especially the black garden ant,
Lasius niger)

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Adaptive Significance

- Selects most profitable from array of food sources
- Selects shortest route to it
 - longer paths abandoned within 1–2 hours
- Adjusts amount of exploration to quality of identified sources
- Collective decision making can be as accurate and effective as some vertebrate individuals

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Observations on Trail Formation

- Two equal-length paths presented at same time: ants choose one at random
- Sometimes the longer path is initially chosen
- Ants may remain “trapped” on longer path, once established
- Or on path to lower quality source, if it’s discovered first
- But there may be advantages to sticking to paths
 - easier to follow
 - easier to protect trail & source
 - safer

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Process of Trail Formation

1. Trail laying
2. Trail following

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Trail Laying

- On discovering food, forager lays chemical trail while returning to nest
 - only ants who have found food deposit pheromone
- Others stimulated to leave nest by:
 - the trail
 - the recruiter exciting nestmates (sometimes)
- In addition to defining trail, pheromone:
 - serves as general orientation signal for ants outside nest
 - serves as arousal signal for ants inside

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Additional Complexities

- Some ants begin marking on return from discovering food
- Others on their first return trip to food
- Others not at all, or variable behavior
- Probability of trail laying decreases with number of trips

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Frequency of Trail Marking

- Ants modulate frequency of trail marking
- May reflect quality of source
 - hence more exploration if source is poor
- May reflect orientation to nest
 - ants keep track of general direction to nest
 - and of general direction to food source
 - trail laying is less intense if the angle to homeward direction is large

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Trail Following

- Ants preferentially follow stronger of two trails
 - show no preference for path they used previously
- Ant may double back, because of:
 - decrease of pheromone concentration
 - unattractive orientation

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Probability of Choosing One of Two Branches

- Let C_L and C_R be units of pheromone deposited on left & right branches
- Let P_L and P_R be probabilities of choosing them
- Then:

$$P_L = \frac{(C_L + 6)^2}{(C_L + 6)^2 + (C_R + 6)^2}$$

- Nonlinearity amplifies probability

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Additional Adaptations

- If a source is crowded, ants may return to nest or explore for other sources
- New food sources are preferred if they are near to existing sources
- Foraging trails may rotate systematically around a nest

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Pheromone Evaporation

- Trails can persist from several hours to several months
- Pheromone has mean lifetime of 30-60 min.
- But remains detectable for many times this
- Long persistence of pheromone prevents switching to shorter trail
- Artificial ant colony systems rely more heavily on evaporation

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