PROBLEM SOLVING AND SEARCH

Chapter 3

What is impact of *sequences* of Percepts?

	Percept sequence	Action
	[A, Clean]	Right
	[A, Dirty]	Suck
	[B, Clean]	Left
\rightarrow	[B, Dirty] [A, Clean], [A, Clean]	Suck
		Right
	[A, Clean], [A, Dirty]	Suck
	•••	

What is difference between agent function and agent program?

Percept sequence	Action
[A, Clean]	Right
[A, Dirty]	Suck
[B, Clean]	Left
[B, Dirty]	Suck
[A, Clean], [A, Clean]	Right
[A, Clean], [A, Dirty]	Suck
•••	

function REFLEX-VACUUM-AGENT([location, status]) returns an action
 if status == Dirty then return Suck
 else if location == A then return Right
 else if location == B then return Left

Outline

- \Diamond Problem-solving agents
- \diamond Problem types
- \diamondsuit Problem formulation
- \diamond Example problems
- \diamond Basic search algorithms

Problem-solving agents

Restricted form of general agent:

```
function SIMPLE-PROBLEM-SOLVING-AGENT( percept) returns an action

static: seq, an action sequence, initially empty

state, some description of the current world state

goal, a goal, initially null

problem, a problem formulation

state \leftarrow UPDATE-STATE(state, percept)

if seq is empty then

goal \leftarrow FORMULATE-GOAL(state)

problem \leftarrow FORMULATE-PROBLEM(state, goal)

seq \leftarrow SEARCH( problem)

action \leftarrow RECOMMENDATION(seq, state)

seq \leftarrow REMAINDER(seq, state)

return action
```

Note: this is offline problem solving; solution executed "eyes closed." Online problem solving involves acting without complete knowledge.

Example: Romania

On holiday in Romania; currently in Arad. Flight leaves tomorrow from Bucharest

Formulate goal:

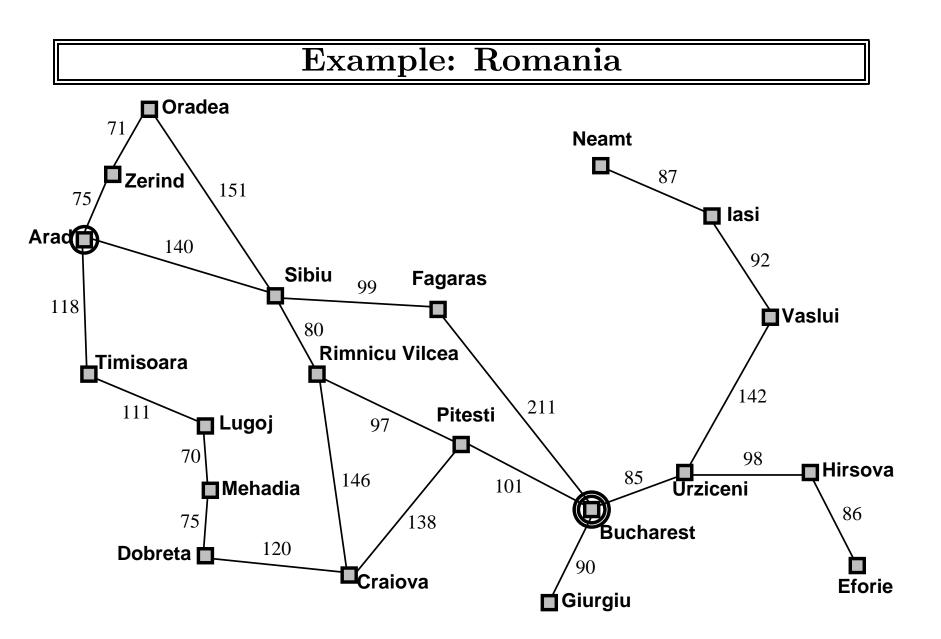
be in Bucharest

Formulate problem:

states: various cities actions: drive between cities

Find solution:

sequence of cities, e.g., Arad, Sibiu, Fagaras, Bucharest



Problem types

Deterministic, fully observable \implies single-state problem Agent knows exactly which state it will be in; solution is a sequence

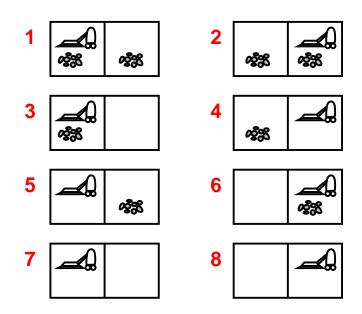
 $\mathsf{Non-observable} \Longrightarrow \mathsf{conformant} \ \mathsf{problem}$

Agent may have no idea where it is; solution (if any) is a sequence

Nondeterministic and/or partially observable \implies contingency problem percepts provide **new** information about current state solution is a contingent plan or a policy often **interleave** search, execution

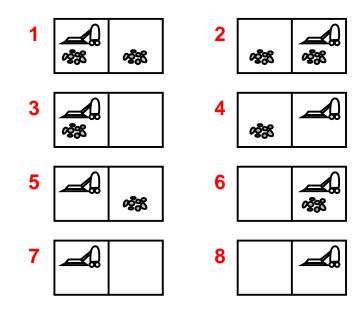
Unknown state space \implies exploration problem ("online")

Single-state, start in #5. <u>Solution</u>??



Single-state, start in #5. <u>Solution</u>?? [*Right*, *Suck*]

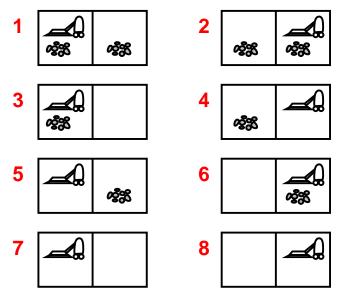
Conformant, start in $\{1, 2, 3, 4, 5, 6, 7, 8\}$ e.g., *Right* goes to $\{2, 4, 6, 8\}$. <u>Solution</u>??



Single-state, start in #5. <u>Solution</u>?? [*Right*, *Suck*]

Conformant, start in $\{1, 2, 3, 4, 5, 6, 7, 8\}$ e.g., *Right* goes to $\{2, 4, 6, 8\}$. <u>Solution</u>?? [*Right*, *Suck*, *Left*, *Suck*]

Contingency, start in #5 Murphy's Law: *Suck* can dirty a clean carpet Local sensing: dirt, location only. <u>Solution</u>??

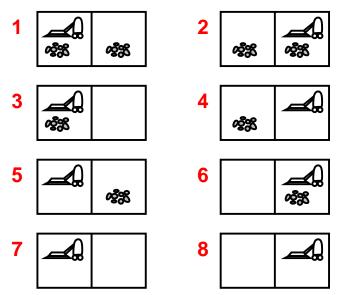


Single-state, start in #5. <u>Solution</u>?? [*Right*, *Suck*]

Conformant, start in $\{1, 2, 3, 4, 5, 6, 7, 8\}$ e.g., *Right* goes to $\{2, 4, 6, 8\}$. <u>Solution</u>?? [*Right*, *Suck*, *Left*, *Suck*]

Contingency, start in #5 Murphy's Law: *Suck* can dirty a clean carpet Local sensing: dirt, location only. <u>Solution</u>??

[Right, if dirt then Suck]



Single-state problem formulation

- A *problem* is defined by five components:
 - Initial state e.g., In(Arad)
 - Actions ACTIONS(s)
 - e.g., {Go(Sibiu), Go(Timisoara), Go(Zerind)}
 - Transition model RESULT(s,a)
 - e.g, RESULT(In(Arad),Go(Zerind)) = In(Zerind)
 - Successor: any state reachable from given state by single action
 - Goal test, can be:
 - Explicit: e.g., x = In(Bucharest)
 - Implicit: e.g., NoDirt(x)
 - Path cost (additive)
 - e.g., sum of distances, number of actions executed, etc.
 - c(s, a, s') is the step cost, assumed to be ≥ 0
- A solution is a sequence of actions leading from initial state to a goal state

Selecting a state space

Real world is absurdly complex

 \Rightarrow state space must be ${\color{black} abstracted}$ for problem solving

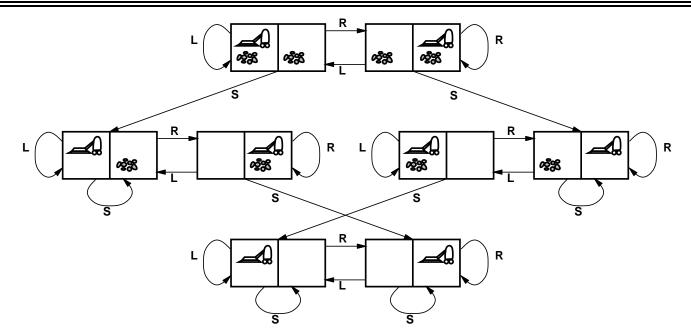
(Abstract) state = set of real states

(Abstract) action = complex combination of real actions e.g., "Arad → Zerind" represents a complex set of possible routes, detours, rest stops, etc.
For guaranteed realizability, any real state "in Arad" must get to some real state "in Zerind"

(Abstract) solution =

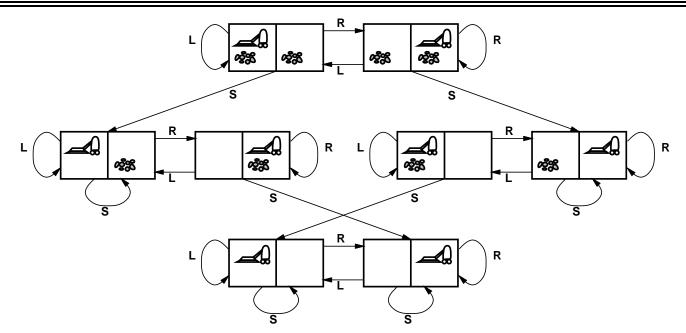
set of real paths that are solutions in the real world

Each abstract action should be "easier" than the original problem!

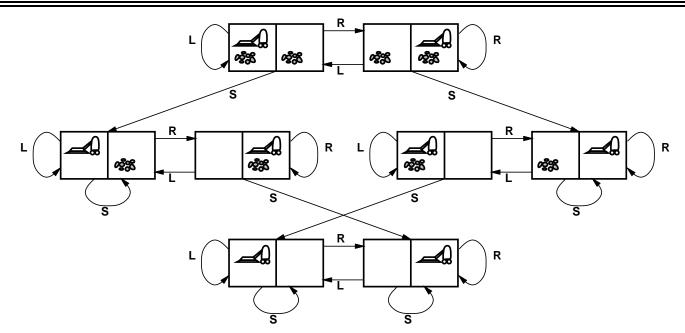


states??
actions??
goal test??
path cost??

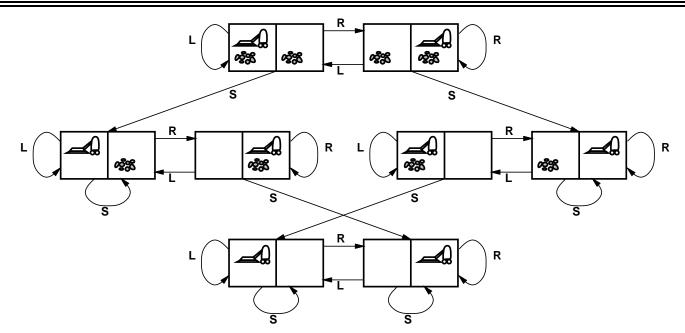
Chapter 3 14



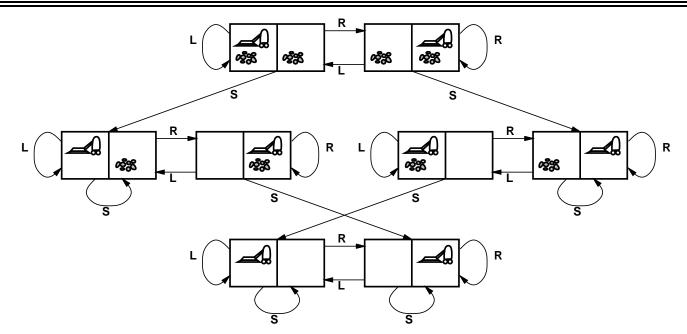
states??: integer dirt and robot locations (ignore dirt amounts etc.)
actions??
goal test??
path cost??



states??: integer dirt and robot locations (ignore dirt amounts etc.)
actions??: Left, Right, Suck, NoOp
goal test??
path cost??

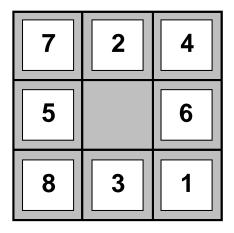


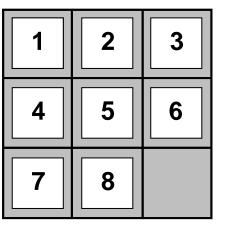
states??: integer dirt and robot locations (ignore dirt amounts etc.)
actions??: Left, Right, Suck, NoOp
goal test??: no dirt
path cost??



states??: integer dirt and robot locations (ignore dirt amounts etc.)
actions??: Left, Right, Suck, NoOp
goal test??: no dirt
path cost??: 1 per action (0 for NoOp)

Example: The 8-puzzle

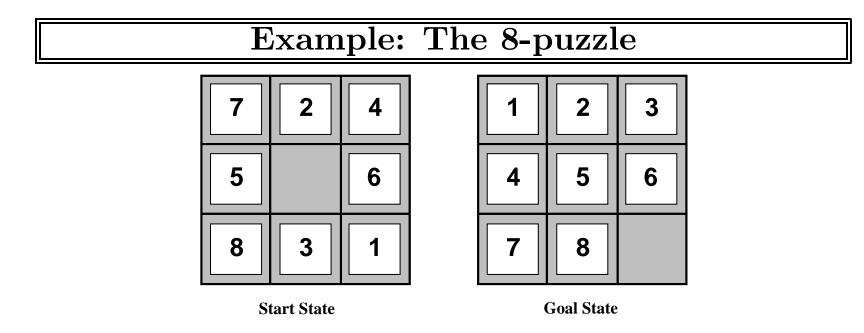




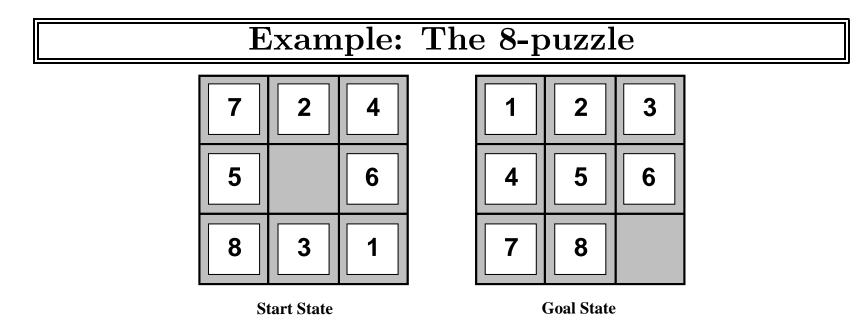
Start State

Goal State

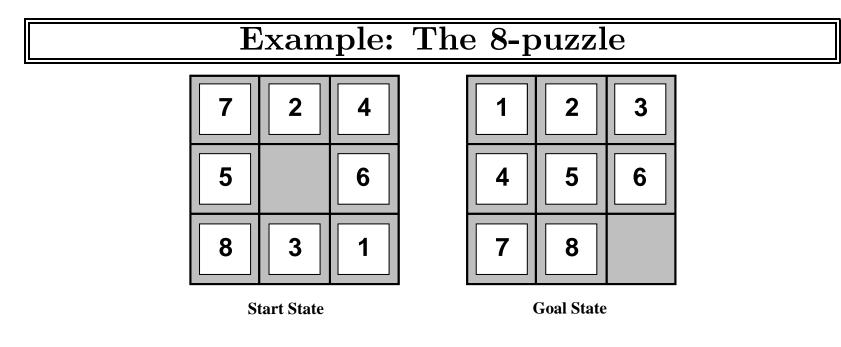
states??
actions??
goal test??
path cost??



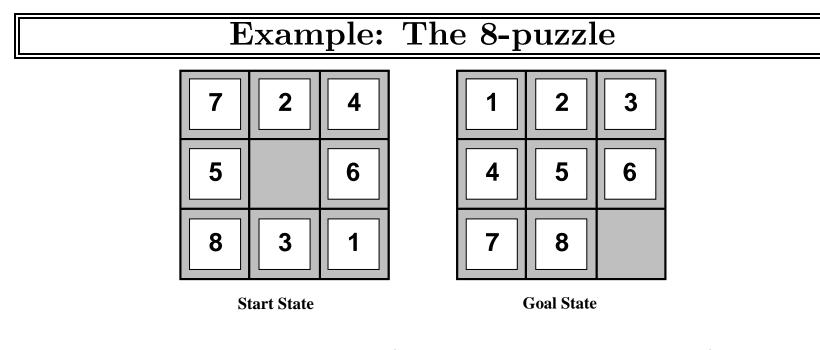
states??: integer locations of tiles (ignore intermediate positions)
actions??
goal test??
path cost??



states??: integer locations of tiles (ignore intermediate positions)
actions??: move blank left, right, up, down (ignore unjamming etc.)
goal test??
path cost??



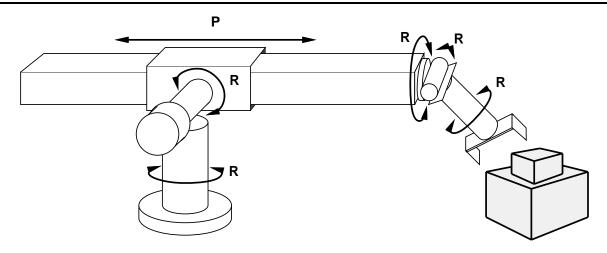
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goal test??: = goal state (given)
path cost??



states??: integer locations of tiles (ignore intermediate positions)
actions??: move blank left, right, up, down (ignore unjamming etc.)
goal test??: = goal state (given)
path cost??: 1 per move

[Note: optimal solution of *n*-Puzzle family is NP-hard]

Example: robotic assembly



states??: real-valued coordinates of robot joint angles parts of the object to be assembled

<u>actions</u>??: continuous motions of robot joints

goal test??: complete assembly with no robot included!

path cost??: time to execute

Exercise – Problem Formulation #1

Give a problem formulation for the following:

1) Using only four colors, you have to color a planar map in such a way that no two adjacent regions have the same color

Initial state?

Actions?

Goal test?

Path cost?

Exercise – Problem Formulation #2

Give a problem formulation for the following:

 A 3-foot-tall monkey is in a room where some bananas are suspended from the 8-foot ceiling. He would like to get the bananas. The room contains two stackable, movable, climbable 3-foot-high crates.

Initial state? Actions?

Goal test?

Path cost?

Exercise – Problem Formulation #3

Give a problem formulation for the following:

3) You have 3 jugs, measuring 12 gallons, 8 gallons, and 3 gallons, and a water faucet. You can fill the jugs up or empty them out from one to another or onto the ground. You need to measure out exactly one gallon.

Initial state?

Actions?

Goal test?

Path cost?

Tree search algorithm

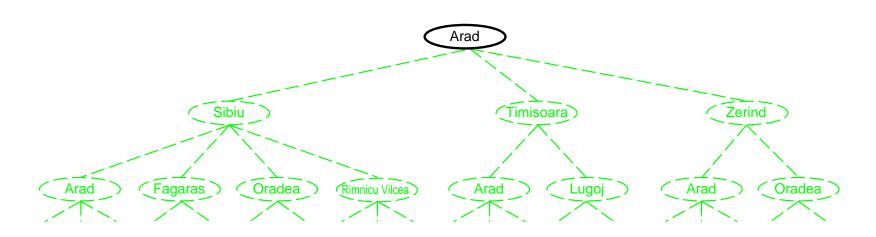
- Basic idea:
 - Offline, simulated exploration of state space by generating successors of already-explored states (i.e., *expanding* states)

function TREE-SEARCH(problem) returns a solution, or failure
 initialize the frontier using the initial state of problem
 loop do
 if the frontier is empty then return failure
 choose a leaf node and remove it from the frontier
 if the node contains a goal state then return the corresponding solution
 expand the chosen node, adding the resulting nodes to the frontier

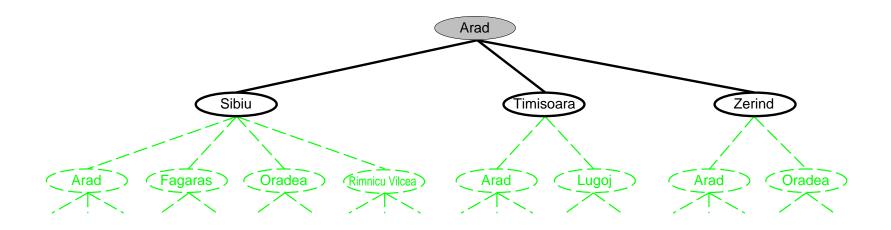
Graph search algorithm

function GRAPH-SEARCH(*problem*) **returns** a solution, or failure initialize the frontier using the initial state of *problem* initialize the explored set to be empty loop do if the frontier is empty then return failure choose a leaf node and remove it from the frontier **if** the node contains a goal state **then return** the corresponding solution add the node to the explored set expand the chosen node, adding the resulting nodes to the frontier only if not in the frontier or explored set

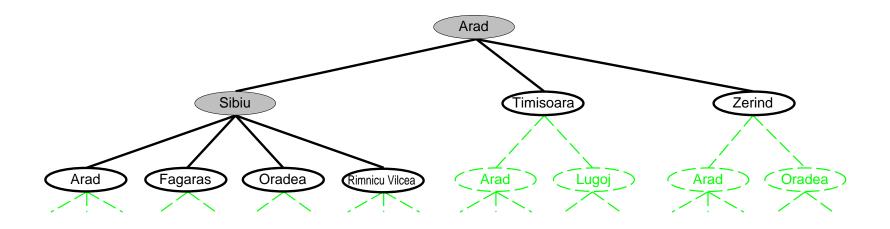
Tree search example



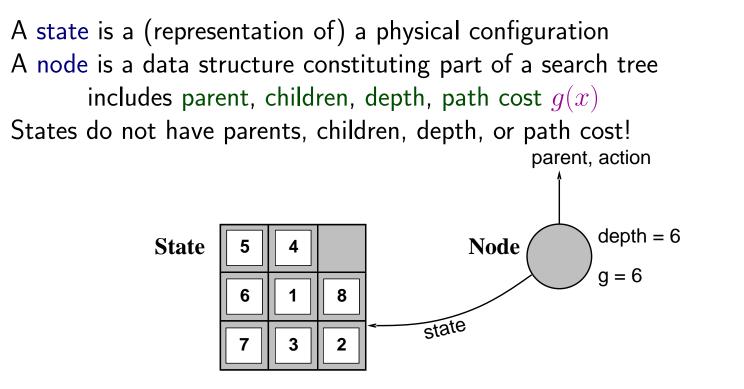
Tree search example



Tree search example



Implementation: states vs. nodes



The Expand function creates new nodes, filling in the various fields and using the SUCCESSORFN of the problem to create the corresponding states.

What's the difference between a world state, a state description, and a search node?

• World state:

• State description:

• Search nodes:

Search strategies

A strategy is defined by picking the order of node expansion

Strategies are evaluated along the following dimensions: completeness—does it always find a solution if one exists? time complexity—number of nodes generated/expanded space complexity—maximum number of nodes in memory optimality—does it always find a least-cost solution?

Time and space complexity are measured in terms of *b*—maximum branching factor of the search tree *d*—depth of the least-cost solution *m*—maximum depth of the state space (may be ∞)

Uninformed search strategies

Uninformed strategies use only the information available in the problem definition

Breadth-first search

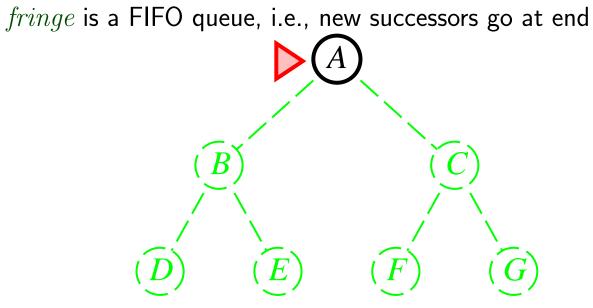
Uniform-cost search

Depth-first search

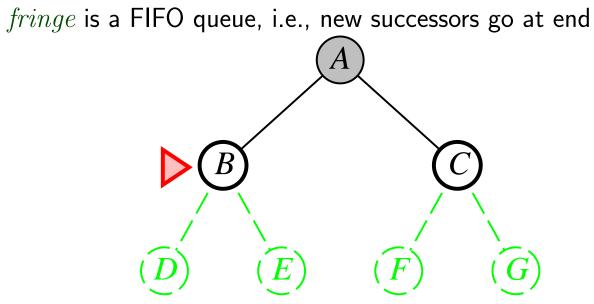
Depth-limited search

Iterative deepening search

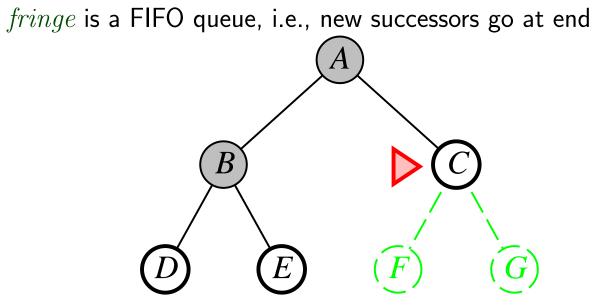
Expand shallowest unexpanded node



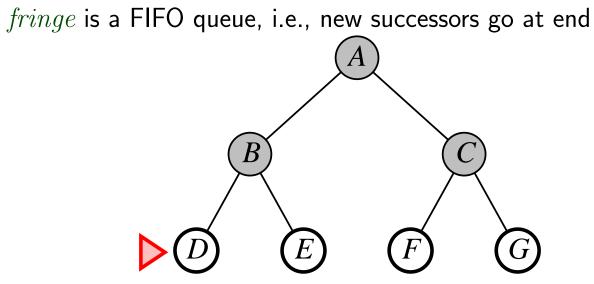
Expand shallowest unexpanded node



Expand shallowest unexpanded node



Expand shallowest unexpanded node



Complete??

<u>Complete</u>?? Yes (if *b* is finite)

Time??

Complete?? Yes (if *b* is finite)

<u>Time</u>?? $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) = O(b^{d+1})$, i.e., exp. in d Space??

Complete?? Yes (if *b* is finite)

<u>Time</u>?? $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) = O(b^{d+1})$, i.e., exp. in d

<u>Space</u>?? $O(b^{d+1})$ (keeps every node in memory)

Optimal??

Complete?? Yes (if *b* is finite)

<u>Time</u>?? $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) = O(b^{d+1})$, i.e., exp. in d

<u>Space</u>?? $O(b^{d+1})$ (keeps every node in memory)

Optimal?? Yes (if cost = 1 per step); not optimal in general

Space is the big problem; can easily generate nodes at 100MB/sec so 24hrs = 8640GB.

Uniform-cost search

Expand least-cost unexpanded node

Implementation:

fringe = queue ordered by path cost, lowest first

Equivalent to breadth-first if step costs all equal

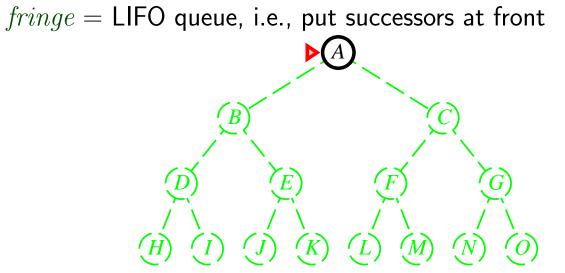
<u>Complete</u>?? Yes, if step cost $\geq \epsilon$

<u>Time</u>?? # of nodes with $g \leq \text{ cost of optimal solution}$, $O(b^{\lceil C^*/\epsilon \rceil})$ where C^* is the cost of the optimal solution

<u>Space</u>?? # of nodes with $g \leq \text{ cost of optimal solution, } O(b^{\lceil C^*/\epsilon \rceil})$

Optimal?? Yes—nodes expanded in increasing order of g(n)

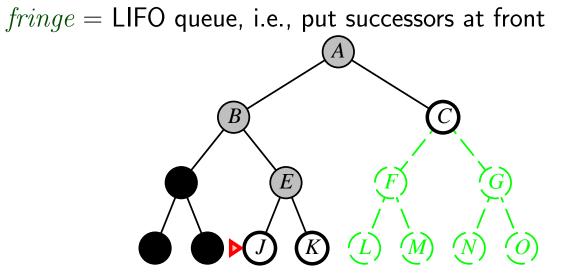
Expand deepest unexpanded node



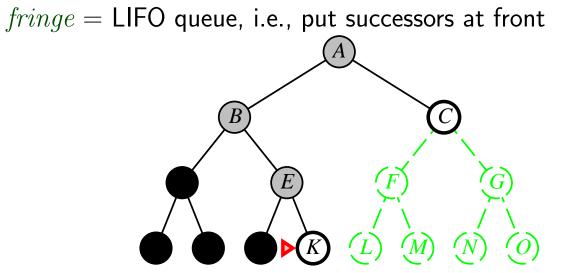
Expand deepest unexpanded node

Implementation:

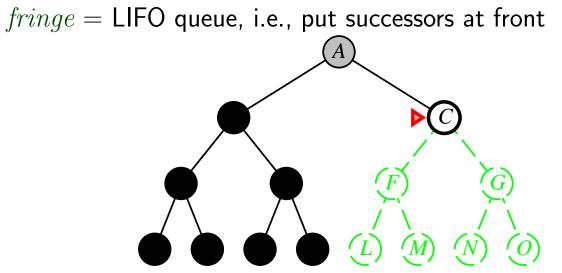
Expand deepest unexpanded node



Expand deepest unexpanded node

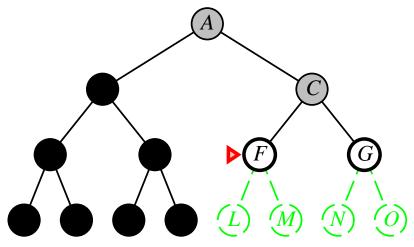


Expand deepest unexpanded node



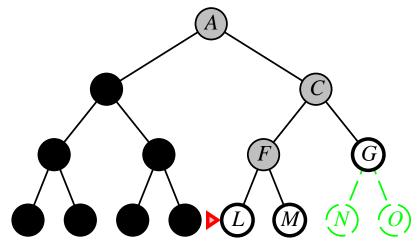
Expand deepest unexpanded node

Implementation:



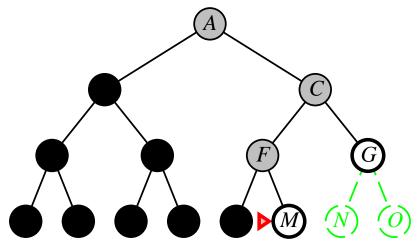
Expand deepest unexpanded node

Implementation:



Expand deepest unexpanded node

Implementation:



Complete??

<u>Complete</u>?? No: fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path ⇒ complete in finite spaces

Time??

<u>Complete</u>?? No: fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path ⇒ complete in finite spaces

<u>Time</u>?? $O(b^m)$: terrible if m is much larger than dbut if solutions are dense, may be much faster than breadth-first

Space??

<u>Complete</u>?? No: fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path ⇒ complete in finite spaces

<u>Time</u>?? $O(b^m)$: terrible if m is much larger than dbut if solutions are dense, may be much faster than breadth-first

Space?? O(bm), i.e., linear space!

Optimal??

<u>Complete</u>?? No: fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path ⇒ complete in finite spaces

<u>Time</u>?? $O(b^m)$: terrible if m is much larger than dbut if solutions are dense, may be much faster than breadth-first

Space?? O(bm), i.e., linear space!

Optimal?? No

Depth-limited search

- Same as depth-first search, but with depth limit *l*
 - i.e., nodes at depth *l* have no successors

```
function DEPTH-LIMITED-SEARCH (problem, limit) returns a solution, or failure/cutoff
return Recursive-DLS(MAKE-NODE(problem, INITIAL-STATE), problem, limit)
```

function RECURSIVE-DLS(node, problem, limit) returns a solution, or failure/cutoff
 if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
 else if limit = 0 then return cutoff

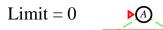
else

cutoff_occurred? ← false
for each action in problem.ACTIONS(node.STATE) do
 child ← Child-Node(problem, node, action)
 result ← Recursive-DLS(child, problem, limit – 1)
 if result = cutoff then cutoff_occurred? ← true
 else if result ≠ failure then return result
 if cutoff_occurred? then return cutoff else return failure

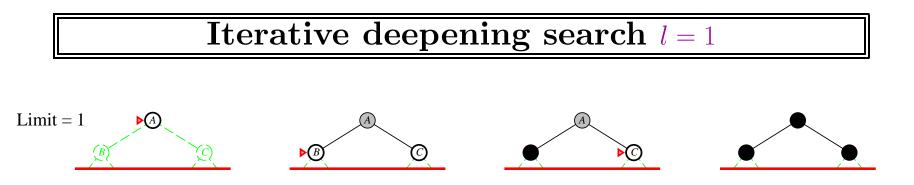
Iterative deepening search

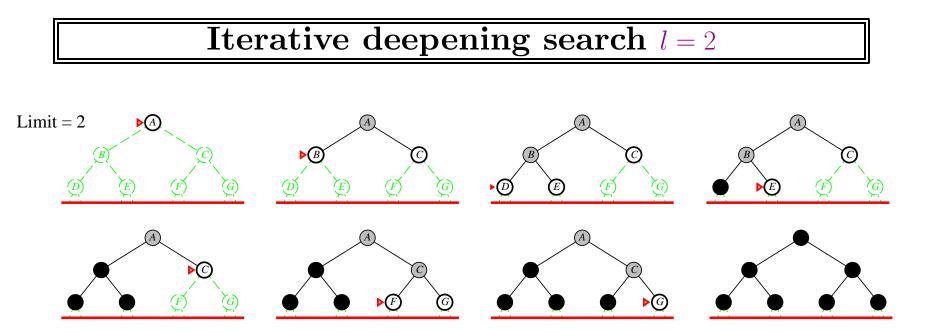
```
function ITERATIVE-DEEPENING-SEARCH( problem) returns a solution
inputs: problem, a problem
for depth ← 0 to ∞ do
    result ← DEPTH-LIMITED-SEARCH( problem, depth)
    if result ≠ cutoff then return result
end
```

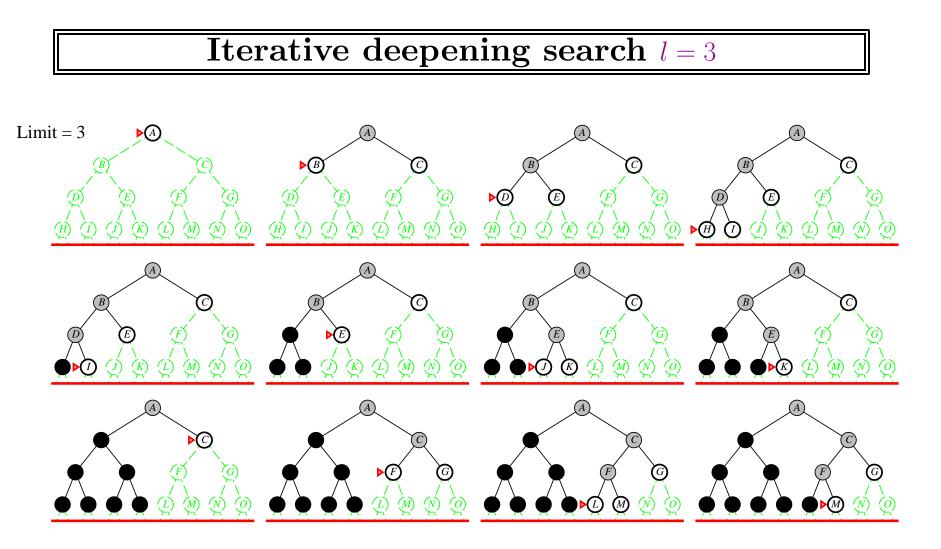
Iterative deepening search l = 0











Properties of iterative deepening search

Complete??

Properties of iterative deepening search

Complete?? Yes

Time??

Properties of iterative deepening search

Complete?? Yes

<u>Time</u>?? $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d = O(b^d)$

Space??

Properties of iterative deepening search

Complete?? Yes

<u>Time</u>?? $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d = O(b^d)$

Space?? O(bd)

Optimal??

Properties of iterative deepening search

Complete?? Yes

<u>Time</u>?? $(d+1)b^0 + db^1 + (d-1)b^2 + \ldots + b^d = O(b^d)$

Space?? O(bd)

Optimal?? Yes, if step cost = 1

Can be modified to explore uniform-cost tree

Numerical comparison for b = 10 and d = 5, solution at far right leaf:

 $N(\mathsf{IDS}) = 50 + 400 + 3,000 + 20,000 + 100,000 = 123,450$ $N(\mathsf{BFS}) = 10 + 100 + 1,000 + 10,000 + 100,000 \times 100,0000 \times 100,000 \times 100,$

HDS does better because ather nodes at depth X are not expanded X

BES can be modified to apply goal test when a hode is generated XX

Summary of algorithms

Criterion	Breadth-	Uniform-	Depth-	Depth-	Iterative
	First	Cost	First	Limited	Deepening
Complete?	Yes^*	Yes ^{*#}	No	Yes, if $l \geq d$ *	Yes *
Time	b^{d+1}	$b^{\lceil C^*/\epsilon \rceil}$	b^m	b^l	b^d
Space	b^{d+1}	$b^{\lceil C^*/\epsilon \rceil}$	bm	bl	bd
Optimal?	Yes***	Yes	No	No	Yes**

* If b is finite

Complete if step costs positive

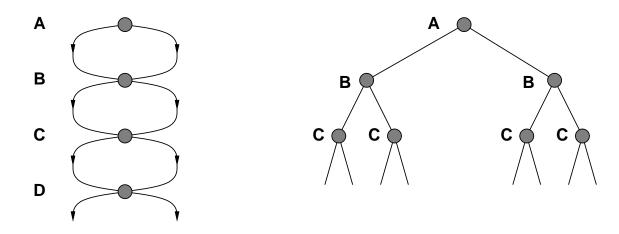
** Optimal if step costs are all identical

Bidirectional search

- Run two simultaneous searches one forward from initial state, and the other backward from the goal
 - We hope they meet in the middle!
- Implemented by changing goal test with a check to see whether the frontiers of the two searches intersect
 - If they do, then a solution is found
- Difficult to use when there are multiple possible goal states
- **Complete?** Yes (if *b* is finite, and both directions use BFS)
- Time: $O(b^{d/2})$
- **Space:** $O(b^{d/2})$
- **Optimal?** Yes (if step costs are identical and both directions use BFS)

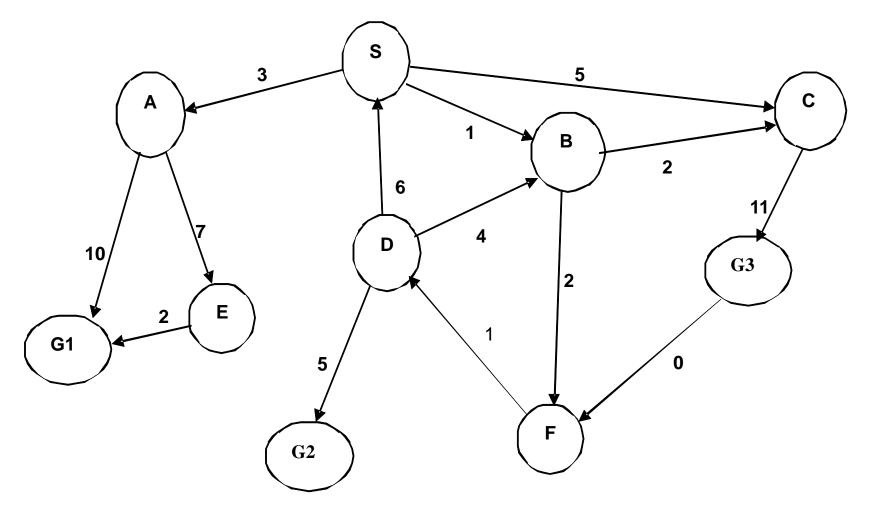
Repeated states

Failure to detect repeated states can turn a linear problem into an exponential one!

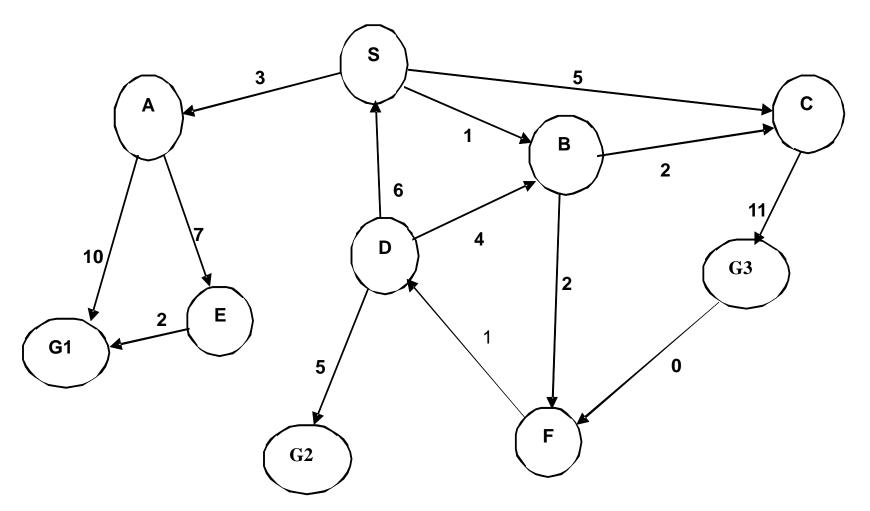


Exercise – Search

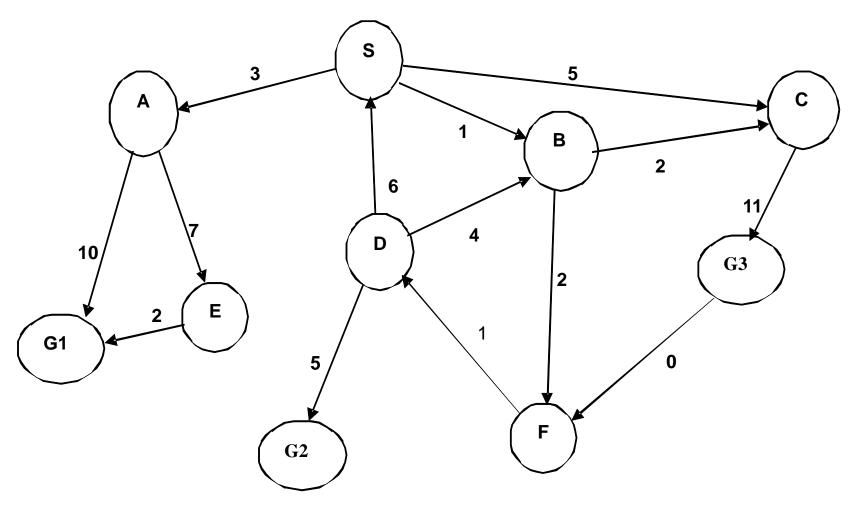
• Consider the search space below, where *S* is the start node and *G1*, *G2*, and *G3* satisfy the goal test. Arcs are labeled with the cost of traversing them. Nodes are removed from fringe in alphabetical order.



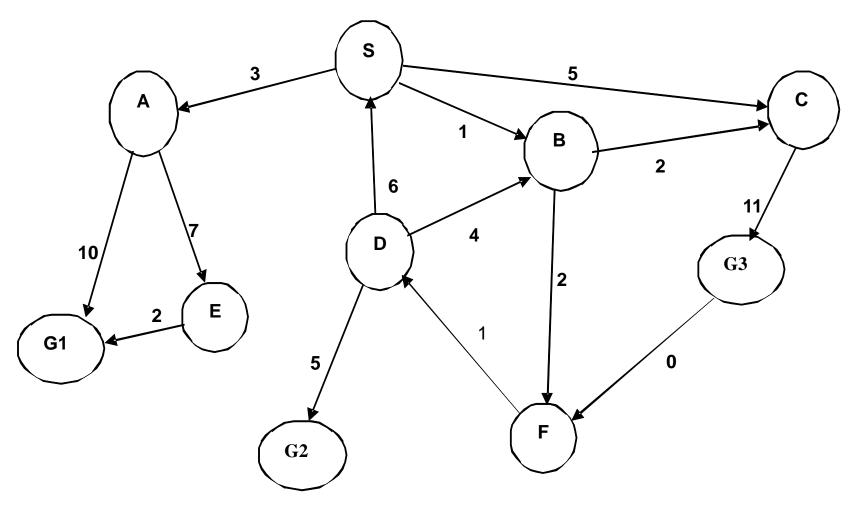
- Iterative Deepening:
 - What is order that nodes are removed from fringe?
 - Which goal state is reached?



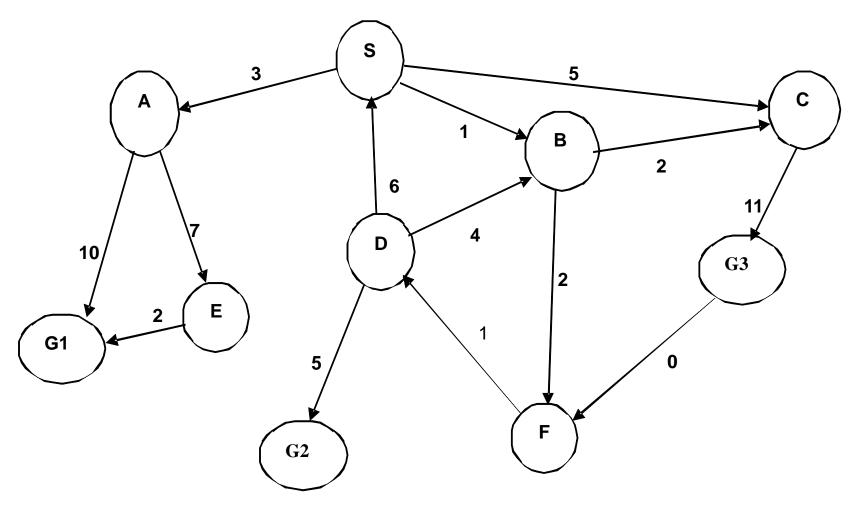
- Iterative Deepening:
 - What is order that nodes are removed from fringe? S
 - Which goal state is reached?



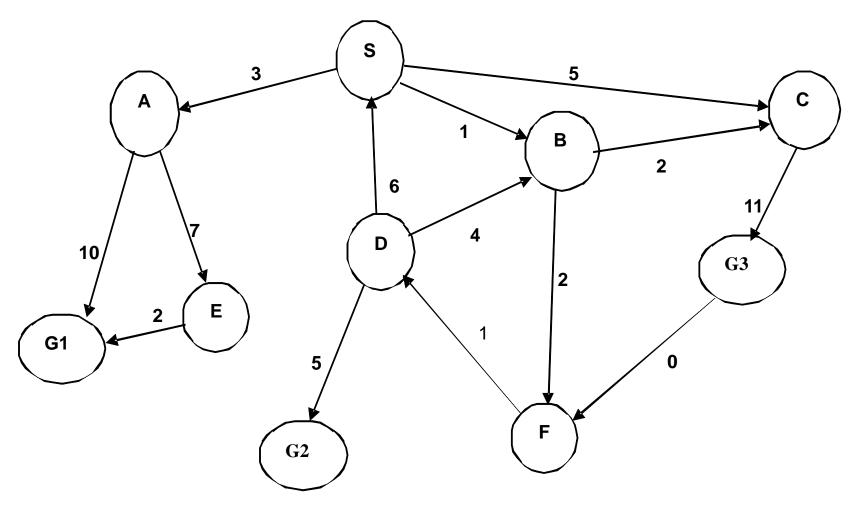
- Iterative Deepening:
 - What is order that nodes are removed from fringe? **S S**
 - Which goal state is reached?



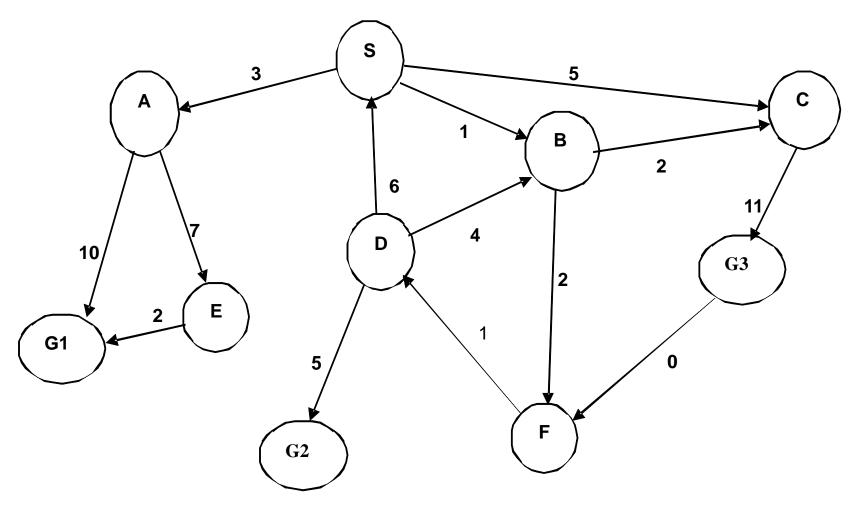
- Iterative Deepening:
 - What is order that nodes are removed from fringe? S S A
 - Which goal state is reached?



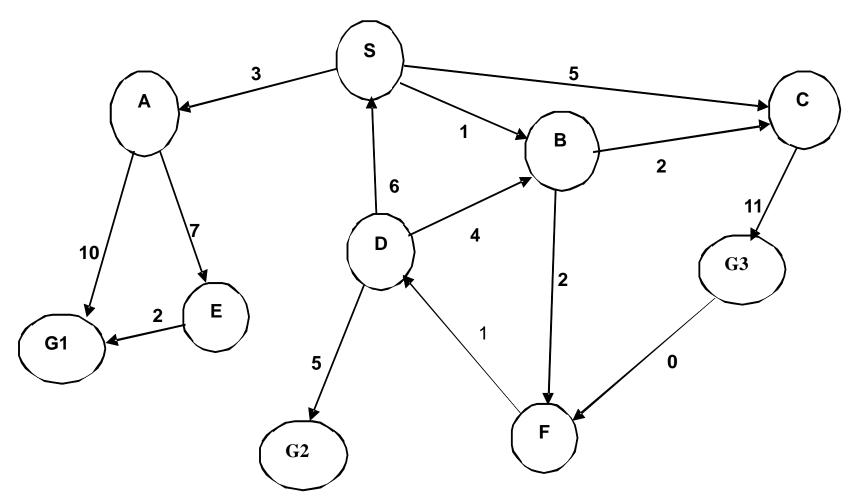
- Iterative Deepening:
 - What is order that nodes are removed from fringe? **S A B**
 - Which goal state is reached?



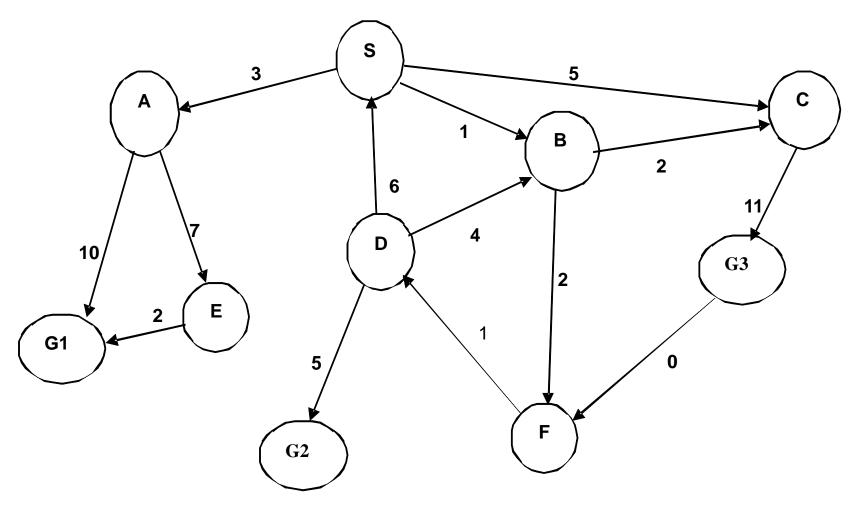
- Iterative Deepening:
 - What is order that nodes are removed from fringe? S S A B C
 - Which goal state is reached?



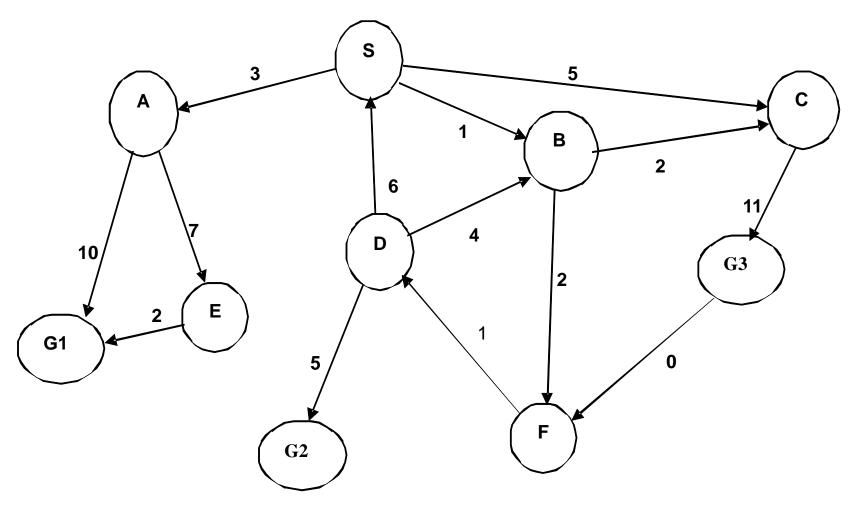
- Iterative Deepening:
 - What is order that nodes are removed from fringe? S S A B C S
 - Which goal state is reached?



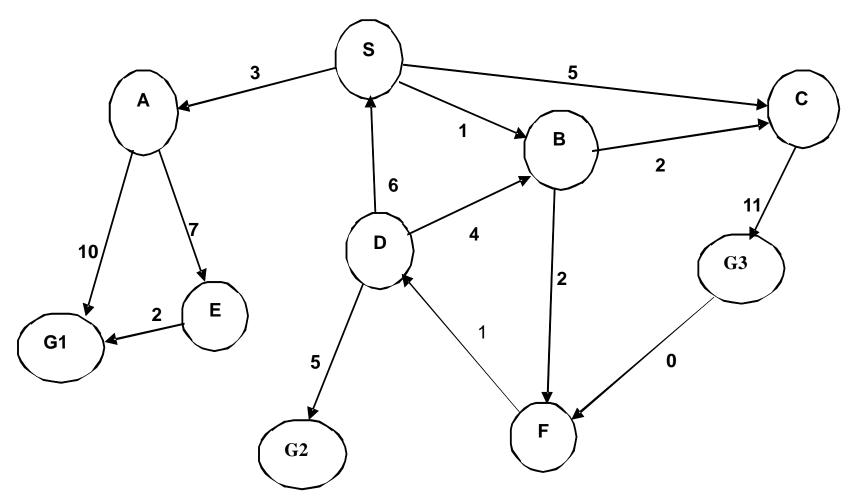
- Iterative Deepening:
 - What is order that nodes are removed from fringe? S S A B C S A
 - Which goal state is reached?



- Iterative Deepening:
 - What is order that nodes are removed from fringe? S S A B C S A E
 - Which goal state is reached?



- Iterative Deepening:
 - What is order that nodes are removed from fringe? **S S A B C S A E G1**
 - Which goal state is reached? G1



Summary

Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored

Variety of uninformed search strategies

Iterative deepening search uses only linear space and not much more time than other uninformed algorithms

Graph search can be exponentially more efficient than tree search

Iterative deepening is the preferred uninformed search method when the search space is large and the depth of the solution is not known.