# Problem solving And search 

Chapter 3

What is impact of sequences of Percepts?


## What is difference between agent function and agent program?

| Percept sequence | Action |
| :--- | :--- |
| $[$ A, Clean $]$ | Right |
| $[$ A, Dirty $]$ | Suck |
| $[B$, Clean $]$ | Left |
| [B, Dirty $]$ | Suck |
| [A, Clean $],$ [A, Clean $]$ | Right |
| [A, Clean], [A, Dirty $]$ | Suck |
| $\ldots$ | $\ldots$ |

```
function REFLEX-VACUUM-AGENT([location, status]) returns an action
    if status == Dirty then return Suck
    else if location ==A then return Right
    else if location == B then return Left
```

$\diamond$ Problem-solving agents
$\diamond$ Problem types
$\diamond$ Problem formulation
$\diamond$ Example problems
$\diamond$ Basic search algorithms

## Problem-solving agents

Restricted form of general agent:
function Simple-Problem-Solving-Agent ( percept) returns an action static: seq, an action sequence, initially empty
state, some description of the current world state goal, a goal, initially null
problem, a problem formulation
state $\leftarrow$ Update-State(state, percept)
if $s e q$ is empty then
goal $\leftarrow$ Formulate-Goal(state)
problem $\leftarrow$ Formulate-Problem(state, goal)
seq $\leftarrow \operatorname{SeARCH}($ problem)
action $\leftarrow$ Recommendation(seq, state)
$s e q \leftarrow \operatorname{REmAINDER}(s e q$, state)
return action

Note: this is offline problem solving; solution executed "eyes closed." Online problem solving involves acting without complete knowledge.

## Example: Romania

On holiday in Romania; currently in Arad.
Flight leaves tomorrow from Bucharest
Formulate goal:
be in Bucharest
Formulate problem:
states: various cities
actions: drive between cities
Find solution:
sequence of cities, e.g., Arad, Sibiu, Fagaras, Bucharest


## Problem types

Deterministic, fully observable $\Longrightarrow$ single-state problem
Agent knows exactly which state it will be in; solution is a sequence
Non-observable $\Longrightarrow$ conformant problem
Agent may have no idea where it is; solution (if any) is a sequence
Nondeterministic and/or partially observable $\Longrightarrow$ contingency problem
percepts provide new information about current state
solution is a contingent plan or a policy often interleave search, execution

Unknown state space $\Longrightarrow$ exploration problem ("online")

## Example: vacuum world

Single-state, start in \#5. Solution??


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Single-state, start in \#5. Solution??
[Right, Suck]
Conformant, start in $\{1,2,3,4,5,6,7,8\}$ e.g., Right goes to $\{2,4,6,8\}$. Solution??


## Example: vacuum world

Single-state, start in \#5. Solution?? [Right, Suck]

Conformant, start in $\{1,2,3,4,5,6,7,8\}$ e.g., Right goes to $\{2,4,6,8\}$. Solution?? [Right, Suck, Left, Suck]

Contingency, start in \#5


Murphy's Law: Suck can dirty a clean carpet Local sensing: dirt, location only. Solution??

## Example: vacuum world

Single-state, start in \#5. Solution?? [Right, Suck]

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Contingency, start in \#5


Murphy's Law: Suck can dirty a clean carpet Local sensing: dirt, location only. Solution??
[Right, if dirt then Suck]

## Single-state problem formulation

- A problem is defined by five components:
- Initial state e.g., In(Arad)
- Actions Actions(s)
- e.g., \{Go(Sibiu), Go(Timisoara), Go(Zerind)\}
- Transition model Result(s,a)
- e.g, Result(In(Arad),Go(Zerind)) $=\operatorname{In}$ (Zerind)
- Successor: any state reachable from given state by single action
- Goal test, can be:
- Explicit: e.g., $x=\ln$ (Bucharest)
- Implicit: e.g., NoDirt(x)
- Path cost (additive)
- e.g., sum of distances, number of actions executed, etc.
- $c\left(s, a, s^{\prime}\right)$ is the step cost, assumed to be $\geq 0$
- A solution is a sequence of actions leading from initial state to a goal state


## Selecting a state space

Real world is absurdly complex
$\Rightarrow$ state space must be abstracted for problem solving
(Abstract) state $=$ set of real states
(Abstract) action $=$ complex combination of real actions
e.g., "Arad $\rightarrow$ Zerind" represents a complex set of possible routes, detours, rest stops, etc.
For guaranteed realizability, any real state "in Arad" must get to some real state "in Zerind"
(Abstract) solution $=$
set of real paths that are solutions in the real world
Each abstract action should be "easier" than the original problem!

## Example: vacuum world state space graph


states??
actions??
goal test??
path cost??

## Example: vacuum world state space graph


states??: integer dirt and robot locations (ignore dirt amounts etc.) actions??
goal test??
path cost??

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actions??: Left, Right, Suck, NoOp
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actions??: Left, Right, Suck, NoOp
goal test??: no dirt
path cost??: 1 per action ( 0 for $N o O p$ )

Example: The 8-puzzle


Start State


Goal State
states??
actions??
goal test??
path cost??

## Example: The 8-puzzle



Start State


Goal State
states??: integer locations of tiles (ignore intermediate positions) actions??
goal test??
path cost??

## Example: The 8-puzzle



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actions??: move blank left, right, up, down (ignore unjamming etc.)
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path cost??

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goal test??: = goal state (given)
path cost??

## Example: The 8-puzzle



Start State


Goal State
states??: integer locations of tiles (ignore intermediate positions)
actions??: move blank left, right, up, down (ignore unjamming etc.)
goal test??: = goal state (given)
path cost??: 1 per move
[Note: optimal solution of $n$-Puzzle family is NP-hard]

states??: real-valued coordinates of robot joint angles
parts of the object to be assembled
actions??: continuous motions of robot joints
goal test??: complete assembly with no robot included!
path cost??: time to execute

## Exercise - Problem Formulation \#1

Give a problem formulation for the following:

1) Using only four colors, you have to color a planar map in such a way that no two adjacent regions have the same color

Initial state?
Actions?
Goal test?
Path cost?

## Exercise - Problem Formulation \#2

Give a problem formulation for the following:
2) A 3-foot-tall monkey is in a room where some bananas are suspended from the 8 -foot ceiling. He would like to get the bananas. The room contains two stackable, movable, climbable 3-foot-high crates.

Initial state?
Actions?

Goal test?
Path cost?

## Exercise - Problem Formulation \#3

Give a problem formulation for the following:
3) You have 3 jugs, measuring 12 gallons, 8 gallons, and 3 gallons, and a water faucet. You can fill the jugs up or empty them out from one to another or onto the ground. You need to measure out exactly one gallon.

Initial state?
Actions?

Goal test?
Path cost?

## Tree search algorithm

- Basic idea:
- Offline, simulated exploration of state space by generating successors of already-explored states (i.e., expanding states)

```
function Tree-SeARCH(problem) returns a solution, or failure
    initialize the frontier using the initial state of problem
    loop do
            if the frontier is empty then return failure
            choose a leaf node and remove it from the frontier
            if the node contains a goal state then return the corresponding solution
            expand the chosen node, adding the resulting nodes to the frontier
```


## Graph search algorithm

function GrAPH-SEARCH(problem) returns a solution, or failure initialize the frontier using the initial state of problem
initialize the explored set to be empty
loop do
if the frontier is empty then return failure
choose a leaf node and remove it from the frontier
if the node contains a goal state
then return the corresponding solution
add the node to the explored set
expand the chosen node, adding the resulting nodes to the frontier only if not in the frontier or explored set




## Implementation: states vs. nodes

A state is a (representation of) a physical configuration
A node is a data structure constituting part of a search tree includes parent, children, depth, path cost $g(x)$
States do not have parents, children, depth, or path cost!


The Expand function creates new nodes, filling in the various fields and using the SUCCESSORFN of the problem to create the corresponding states.

# What's the difference between a world state, a state description, and a search node? 

- World state:
- State description:
- Search nodes:


## Search strategies

A strategy is defined by picking the order of node expansion
Strategies are evaluated along the following dimensions: completeness-does it always find a solution if one exists? time complexity-number of nodes generated/expanded space complexity-maximum number of nodes in memory optimality—does it always find a least-cost solution?

Time and space complexity are measured in terms of
$b$-maximum branching factor of the search tree
$d$-depth of the least-cost solution
$m$-maximum depth of the state space (may be $\infty$ )

## Uninformed search strategies

Uninformed strategies use only the information available in the problem definition

Breadth-first search
Uniform-cost search
Depth-first search
Depth-limited search
Iterative deepening search

## Breadth-first search

## Expand shallowest unexpanded node

## Implementation:

fringe is a FIFO queue, i.e., new successors go at end


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Properties of breadth-first search
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Optimal??

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Complete?? Yes (if $b$ is finite)
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Space?? $O\left(b^{d+1}\right)$ (keeps every node in memory)
Optimal?? Yes (if cost $=1$ per step); not optimal in general
Space is the big problem; can easily generate nodes at $100 \mathrm{MB} / \mathrm{sec}$ so $24 \mathrm{hrs}=8640 \mathrm{~GB}$.

## Uniform-cost search

## Expand least-cost unexpanded node

## Implementation:

fringe $=$ queue ordered by path cost, lowest first
Equivalent to breadth-first if step costs all equal
Complete?? Yes, if step cost $\geq \epsilon$
Time?? \# of nodes with $g \leq$ cost of optimal solution, $O\left(b^{\left\lceil C^{*} / \epsilon\right\rceil}\right)$ where $C^{*}$ is the cost of the optimal solution

Space?? \# of nodes with $g \leq$ cost of optimal solution, $O\left(b^{\left[C^{*} / \epsilon\right\rceil}\right)$
Optimal?? Yes—nodes expanded in increasing order of $g(n)$

## Depth-first search

Expand deepest unexpanded node
Implementation:
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Modify to avoid repeated states along path
$\Rightarrow$ complete in finite spaces
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Space??

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Space?? $O(b m)$, i.e., linear space!
Optimal??

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Modify to avoid repeated states along path
$\Rightarrow$ complete in finite spaces
Time?? $O\left(b^{m}\right)$ : terrible if $m$ is much larger than $d$ but if solutions are dense, may be much faster than breadth-first

Space?? $O(b m)$, i.e., linear space!
Optimal?? No

## Depth-limited search

- Same as depth-first search, but with depth limit $l$
- i.e., nodes at depth $l$ have no successors
function Depth-Limited-Search (problem, limit) returns a solution, or failure/cutoff return Recursive-DLS(MAKe-Node(problem, Initial-State), problem, limit)
function ReCURSIVE-DLS(node, problem, limit) returns a solution, or failure/cutoff if problem. Goal-TESt(node.State) then return Solution(node) else if limit $=0$ then return cutoff else
cutoff_occurred? $\leftarrow$ false for each action in problem.Actions(node.State) do child $\leftarrow$ Child-Node(problem, node, action) result $\leftarrow$ Recursive-DLS(child, problem, limit - 1) if result $=$ cutoff then cutoff_occurred? $\leftarrow$ true else if result $\neq$ failure then return result if cutoff_occurred? then return cutoff else return failure


## Iterative deepening search

function ITERATIVE-DEEPENING-SEARCH( problem) returns a solution inputs: problem, a problem
for depth $\leftarrow 0$ to $\infty$ do
result $\leftarrow$ DEPTH-Limited-SEARCH $($ problem, depth $)$
if result $\neq$ cutoff then return result
end

Iterative deepening search $l=0$

Limit $=0$ 도

Iterative deepening search $l=1$


## Iterative deepening search $l=2$



## Iterative deepening search $l=3$







Properties of iterative deepening search
Complete??

## Properties of iterative deepening search

## Complete?? Yes

Time??

## Properties of iterative deepening search

## Complete?? Yes

Time? ? $(d+1) b^{0}+d b^{1}+(d-1) b^{2}+\ldots+b^{d}=O\left(b^{d}\right)$
Space??

## Properties of iterative deepening search

## Complete?? Yes

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Space?? $O(b d)$
Optimal??

## Properties of iterative deepening search

## Complete?? Yes

Time? ? $(d+1) b^{0}+d b^{1}+(d-1) b^{2}+\ldots+b^{d}=O\left(b^{d}\right)$
Space?? $O(b d)$
Optimal?? Yes, if step cost $=1$
Can be modified to explore uniform-cost tree
Numerical comparison for $b=10$ and $d=5$, solution at far right leaf:

$$
\begin{aligned}
N(\mathrm{IDS})=50+400+3,000+20,000+100,000 & =123,450 \\
N(\mathrm{BFS})=10+100+1,000+10,000+100,000 & =111,110
\end{aligned}
$$




## Summary of algorithms

| Criterion | Breadth- <br> First | Uniform- <br> Cost | Depth- <br> First | Depth- <br> Limited | Iterative <br> Deepening |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Complete? | Yes* $^{*}$ | Yes** | No | Yes, if $l \geq d^{*}$ | Yes* |
| Time | $b^{d+1}$ | $b^{\left\lceil C^{*} / \epsilon\right\rceil}$ | $b^{m}$ | $b^{l}$ | $b^{d}$ |
| Space | $b^{d+1}$ | $b^{\left\lceil C^{*} / \epsilon\right\rceil}$ | $b m$ | $b l$ | $b d$ |
| Optimal? | Yes** $^{* *}$ | Yes | No | No | Yes** |

* If $b$ is finite
\# Complete if step costs positive
** Optimal if step costs are all identical


## Bidirectional search

- Run two simultaneous searches - one forward from initial state, and the other backward from the goal
- We hope they meet in the middle!
- Implemented by changing goal test with a check to see whether the frontiers of the two searches intersect
- If they do, then a solution is found
- Difficult to use when there are multiple possible goal states
- Complete? Yes (if $b$ is finite, and both directions use BFS)
- Time: $O\left(b^{d / 2}\right)$
- Space: $O\left(b^{d / 2}\right)$
- Optimal? Yes (if step costs are identical and both directions use BFS)


## Repeated states

Failure to detect repeated states can turn a linear problem into an exponential one!


## Exercise - Search

- Consider the search space below, where $S$ is the start node and G1, G2, and $G 3$ satisfy the goal test. Arcs are labeled with the cost of traversing them. Nodes are removed from fringe in alphabetical order.



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe?
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? $S$
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S A
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S SAB
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S A B C
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S B C
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S B C S A
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S SABCSAE
- Which goal state is reached?



## Exercise - Search (con't)

- Iterative Deepening:
- What is order that nodes are removed from fringe? S S A B S A E G1
- Which goal state is reached? G1



## Summary

Problem formulation usually requires abstracting away real-world details to define a state space that can feasibly be explored

Variety of uninformed search strategies
Iterative deepening search uses only linear space and not much more time than other uninformed algorithms

Graph search can be exponentially more efficient than tree search

Iterative deepening is the preferred uninformed search method when the search space is large and the depth of the solution is not known.

