Topcoder SRM 641, D1, 250-Pointer "TrianglesContainOrigin"

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The problem

• You are given the (x,y) values of points on a twodimensional grid:



The problem

 Of all the triangles with these points as endpoints, how many have the origin inside?





• Example 1: There are 4 triangles, two of which include the origin.





Prototype and Constraints

- **Class name:** TrianglesContainOrigin
- Method: count()
- Parameters:

X	vector <int></int>	X coordinates
У	vector <int></int>	Y coordinates

- Return Value: long long
- Constraints:
 - x.size() == y.size() ≤ 2500 .
 - \times and γ values between -10,000 and 10,000.
 - No three values co-linear
 - No two values co-linear with the origin.

Brain-dead enumeration of triangles

- Let n = x.size().
- Then the number of triangles is:

$$\left(\begin{array}{c}n\\3\end{array}\right) = O(n^3)$$

- When *n* = 2500, this is 2,590,630,000. Too slow.
- Our solution can be $O(n^2)$, but not much slower.

The Key Insight

• Draw a line from the origin to each point, and then calculate the angles of adjacent lines:



The Key Insight

• Consider a triangle – it will include the origin if and only if each of the angles of lines from the origin is less than 180:





The Strategy

- Enumerate all pairs of points whose angle to the origin is less than 180 degrees:
- For a given pair, there is a minimum and maximum angle that the third point can have.



An Algorithm

- For each point, calculate the point's angle *a* from the origin.
- Insert the points into a map keyed by angle.
- Also insert the points keyed by angle+360.
- Number the points in the map by ascending angle.

Let's look at example 3



An Algorithm, Continued

- For each point *x* < 360 and each point *y* whose angle to *x* is less than 180 degrees, use *upper_bound()* to find:
 - The smallest point whose angle is > 180 to *x*.
 - The smallest point whose angle is > 180 to y.
 - The difference in vals is the number of points that can complete the triangle!
- Sum the triangles and divide by 3 for the answer!

Let's look at example 3



Consider points 0 and 4.

Upper-bound will find points 9 and 13.

Therefore, there are four points that can complete the triangle with points 0 and 4.

Improvements

- Insert a giant sentinel, and:
- You only have to insert each point once.
- You don't have to divide by three.
- (0 triangles when b-a > 180)



Time for some clicker questions

We can make it faster.

Label the points: P1, P2, P3, P4



We can make it faster.

Label the points: P1, P2, P3, P4.

- The number of triangles for (P1,P2) is (P4->val P3->val)
- Set P3, and set P4 = P3 when you set P1.
- For each P2, increment P4 until it's right.
- A total of O(n) increments for all P2 makes it $O(n^2)$ overall.



Experiment

- MacBook Pro 2.2 Ghz, optimized with -O2
- Increments of 100, each point the average of ten runs.



How did the Topcoders Do?

- 580 competitors
- 285 (48%) submitted a solution.
- 216 (76%) of the submissions were correct.
- That's 38% I suspected this one would be hard!

(For class)

- Go over the program that makes the jgraph (in the lecture notes directory).
- But delete the part that does the calculation, so you can do the calculation live (maybe even give the students a chance to do it).
- Remember txt/points-500.txt and txt/points-2500.txt.

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