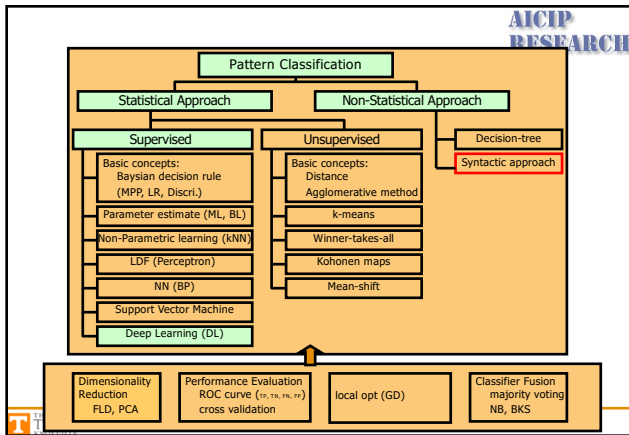



ECE471-571 – Pattern Recognition

Lecture 18: Syntactic Pattern Recognition

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References

- King Sun Fu, *Syntactic Pattern Recognition, Applications*. Springer, 1977
- International Association for Pattern Recognition (IAPR), 1976
- TPAMI, 1978

– DL Davies, DW Bouldin - Pattern Analysis and Machine ..., 1979 (citation: 2275)

3

Key Concept

- ◆ If we can draw it (automatically), then we can recognize it
- ◆ Based on formal language

Philosophy

- ◆ A **grammar** generates a (possibly infinite) set of strings (pictures)
- ◆ If we can design a grammar which generates a class of strings, then we can build a machine which will recognize any string in that class

Types of Grammars - Symbols

- ◆ V_N : the set of non-terminal symbols
- ◆ V_T : the set of terminal symbols
- ◆ P : the set of rewriting rules (productions)
- ◆ S : the start symbol
- ◆ \emptyset : the empty (null) symbol

Type 0 Grammar

- ◆ No restrictions on rewriting rules

$$\alpha \rightarrow \beta$$

- ◆ The string α (whenever it occurs in a derivation) may be replaced by the string β

Type 1 – Context Sensitive

$\alpha \rightarrow \beta$ is allowed only if $|\alpha| \leq |\beta|$ (length)

Type 2 – Context Free

- ◆ Left side must be a single non-terminal

$$A \rightarrow \alpha$$

- ◆ Example

$$S \rightarrow 0S1$$

$$S \rightarrow 01$$

Type 3 - Regular

- ◆ $A \rightarrow aB$, or $A \rightarrow a$
- ◆ A and B are single non-terminal
- ◆ Is a regular grammar also context-free?

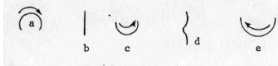
Example



- ◆ Describe two types of chromosomes for recognition (submedian chromosome and telocentric chromosome)
- ◆ Chromosome is represented as a string, obtained by tracing the outline in clockwise direction
- ◆ Pattern primitives = terminal symbols



Example (cont')



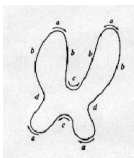
- Grammar for recognition of submedian and telocentric chromosomes
 - $G = (V_N, V_T, P, S)$
- Non-terminals
 - $V_N = \{S, S_1^*, S_2^*, A, B, C, D, E, F\}$
 - S – start symbol
 - S_1^* – submedian chromosome
 - S_2^* – telocentric chromosome
 - A – armpair, B – bottom, C – side, D – arm, E – rightpart, F – leftpart

Example (cont')

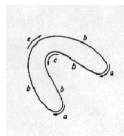
◆ Production (rewriting rules)

$S \rightarrow S_1^*$	$B \rightarrow e$	$S \rightarrow S_2^*$	$C \rightarrow bC$
$S_1^* \rightarrow AA$	$C \rightarrow Cb$	$S_2^* \rightarrow BA$	$C \rightarrow b$
$A \rightarrow CA$	$C \rightarrow d$	$A \rightarrow AC$	$D \rightarrow bD$
$A \rightarrow DE$	$D \rightarrow Db$	$A \rightarrow FD$	$D \rightarrow a$
$B \rightarrow bD$	$E \rightarrow cD$	$B \rightarrow Bb$	$F \rightarrow Dc$

Example (cont')



babcbabdacad



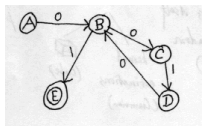
ebabcbab

$S \rightarrow S_1^* \rightarrow AA \rightarrow ACA \rightarrow FDCA \rightarrow DcDCA \rightarrow$
 $bDcDCA \rightarrow bDbcDCA \rightarrow babcDCA \rightarrow$
 $babcDCA \rightarrow babcDbCA \rightarrow babcbabCA \rightarrow$
 $babcabdA \rightarrow babcabdAC \rightarrow babcabdDEC \rightarrow$
 $babcabdEC \rightarrow babcabdacDC \rightarrow$
 $babcabdacaC \rightarrow babcabdacad$

Finite State Machine

◆ A regular expression determines a finite-state machine

■ $0(010)^*1$



■ $S \rightarrow A, A \rightarrow 0B, B \rightarrow 0C, C \rightarrow 1D, D \rightarrow 0B, B \rightarrow 1$

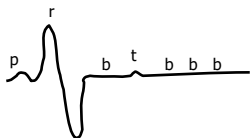
Recognition of Abnormal ECG

◆ Regular grammar

■ $G = (\{S, A, B, C, D, E, H\}, \{p, r, t, b\}, P, S)$

◆ Productions:

■ $S \rightarrow pA, A \rightarrow rB, B \rightarrow bC, C \rightarrow tD, D \rightarrow b, D \rightarrow bE, E \rightarrow b, E \rightarrow bH, E \rightarrow pA, H \rightarrow b, H \rightarrow bS, H \rightarrow pA$



ECG (cont')

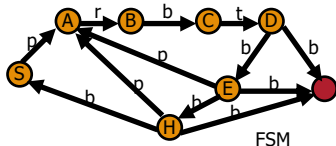
• Example of derivation of a well formed ECG wave:

wave:

– $S \rightarrow pA \rightarrow prB \rightarrow prbC \rightarrow prbtD \rightarrow prbtbE \rightarrow prbtbbH \rightarrow prbtbbbS \rightarrow prbtbbb pA \rightarrow prbtbbb prB \rightarrow prbtbbb prbC \rightarrow prbtbbb prbtD \rightarrow prbtbbb prbtbE \rightarrow prbtbbb prbtbb \rightarrow \dots$ etc.

- Note possibility of variable number of “b’ s”
 - One to three to accommodate normal variations of heart rate

The FSM



Education is what remains
after one has forgotten
everything one learned in
school. -- Albert Einstein
