

## Problem Set 12: *Computational Complexity*

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**Due: Thursday, April 24, 2014, at the beginning of class**

0. Complete the course SAIS evaluation for this course, and turn in your confirmation sheet. I value your constructive feedback!
1. Consider again the 0-1 Knapsack problem: A thief robbing a store finds  $n$  items. The  $i$ th item is worth  $v_i$  dollars and weighs  $w_i$  pounds, where  $v_i$  and  $w_i$  are integers. The thief wants to take as valuable a load as possible, but s/he can carry at most  $W$  pounds in her/his knapsack, for some integer  $W$ . There is dynamic programming solution to this problem that runs in  $O(n W)$  time. [An aside: a useful study exercise for the final is to develop this dynamic programming solution. But, you don't have to show it for this problem set.]  
For this homework, answer this question: Is this dynamic programming solution (which runs in  $O(n W)$  time) a polynomial-time algorithm? Explain your answer.
2. Show that an otherwise polynomial-time algorithm that makes at most a constant number of calls to polynomial-time subroutines runs in polynomial time, but that a polynomial number of calls to polynomial-time subroutines may result in an exponential-time algorithm.
3. Show that if HAM-CYCLE  $\in P$ , then the problem of printing (in order) the vertices of a Hamiltonian cycle is polynomial-time solvable. In other words, give a polynomial-time algorithm that prints out (in order) the vertices of a Hamiltonian cycle in a graph using the assumed polynomial-time subroutine that decides HAM-CYCLE.
4. Let  $5\text{-CLIQUE} = \{\langle G \rangle \mid G \text{ is an undirected graph having a complete subgraph with 5 nodes}\}$ . Show that  $5\text{-CLIQUE}$  is in  $P$ .
5. We define a monotone Boolean formula as a formula with no negated variables. We further define the 2-monotone-small-SAT problem is as follows: Given a monotone Boolean formula  $\Phi$  in 2-CNF form (yes, that's two-CNF) and an integer  $k$ , determine if  $\Phi$  has a truth assignment with no more than  $k$  variables assigned "True" (i.e., "1"). Show that this problem is NP-complete. [Hint: Use Vertex Cover]