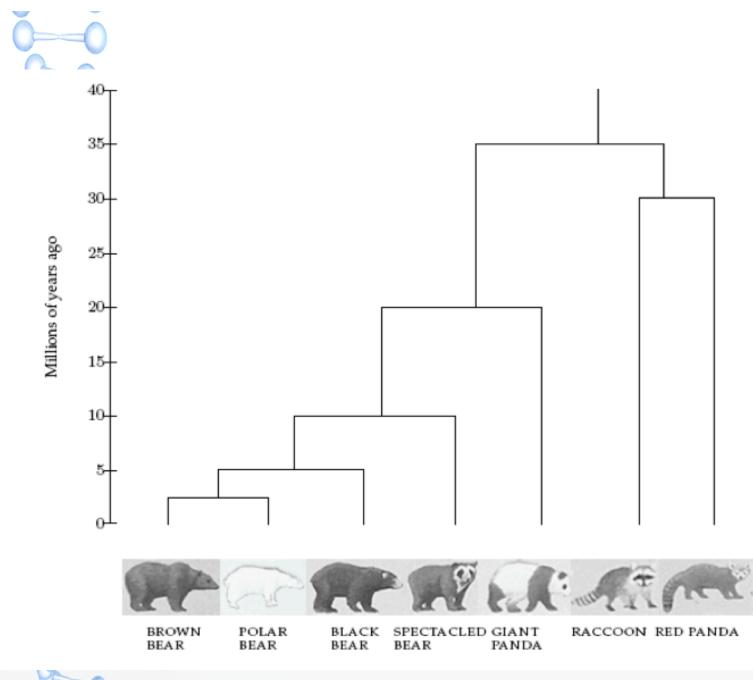
Other approaches to inferring phylogenetic trees



Review

UPGMA

This method requires ultrametric property: for three distances, two are equal and third is <= the first two Relies on a molecular clock

Neighbor joining

This method requires additive property: distances between two nodes is sum of their edges Often produces a good tree even with non-additive

data

JUPGMA Method

Unweighted Pair Group Method using Arithmetic Averages

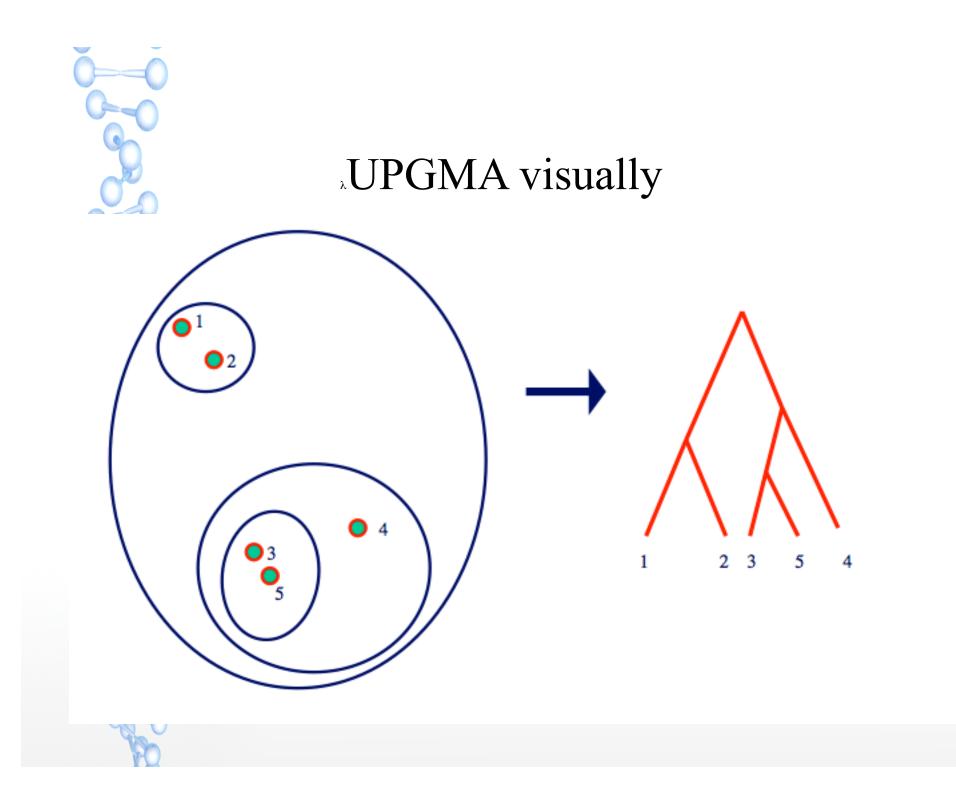
Distance is defined between two clusters *Ci* and *Cj* such that:

 $d_{ij} = \frac{1}{|C_i||C_j|} \sum_{p \in C_i, q \in C_j} d_{pq}$

Basic idea

Dij is the average distance between pairs of taxa from each cluster

Algorithm: Start with one taxa per cluster Iteratively pick two clusters and merge Create a new node in the tree for the merged cluster



More specifics

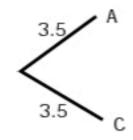
Place each taxon at height 0 in the tree

While more than two clusters: Determine clusters with smallest *dij* Merge clusters into a new one *Ck* Make a new node *k* at height *dij*/2 Replace *Ci* and *Cj* with *Ck* Recompute distance of *Ck* to other clusters

Hook in the two remaining clusters to the root with height calculated as above.

Step 1: pick smallest and update distances

	A	B	С	D	
A	0				
B	8	0			
С	7	9	0		
D	12	14	11	0	

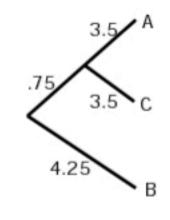


 $\begin{array}{l} M_{B(AC)} = (M_{BA} + M_{BC})/2 = (8+9)/2 = 8.5 \\ M_{D(AC)} = (M_{DA} + M_{DC})/2 = (12+11)/2 = 11.5 \end{array}$

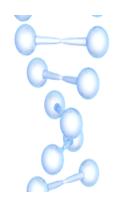




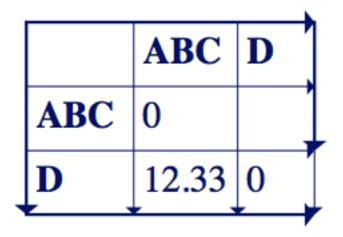
	AC	B	D	
AC	0			
B	8.5	0		,
D	11.5	14	0	_

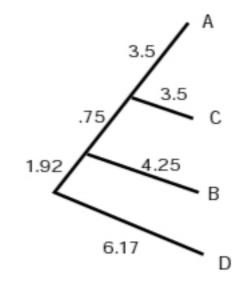


 $M_{(ABC)D} = (M_{AD} + M_{BD} + M_{CD})/3 = (12+14+11)/3$



»Step 3







Molecular clocks

- A molecular clock assumption is divergence is uniform and equal across all branches of the tree
- Seldom (never?) true in practice
- If it is true, these data are called *ultrametric*.

»Neighbor joining

Does not assume a molecular clock, but does assume additively

Distance between a pair of leaves is sum of edges between them

Constructs an unrooted tree iteratively, just like UPGMA

Two differences:

How subtrees selected

How distances are updated

Root can be added via inclusion of an "outgroup"



Basics

- NJ is a greedy algorithm that starts with a center star tree (all taxa connected to a single root)
- Criterion for merging is key: identifies topological neighbors using math that is correct for all additive distance matrices.
- Once merged, the two taxa are treated as a single taxon.

Idea behind NJ algorithm

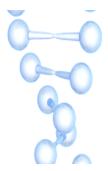
Start with a center star phylogeny where everyone is connected to a single node.

Choose two sequences to merge based on the mathematically optimal topology under an additive scoring scheme.

Note this tree starts out unrooted and remains so without an outgroup.

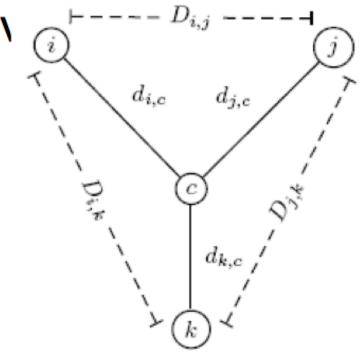
Caveats

- Matrix is updated iteratively after merge.
- Produces unrooted trees; need an outgroup for a rooted version
- Always gives the true tree if distances are additive (may not with noise)



Reconstructing a 3 leaved tree

- Tree reconstruction for any 3x3 matrix is straightforward
- We have 3 leaves *i*, *j*, *k* and a center



<u>Observe:</u>

$$d_{ic} + d_{jc} = D_{ij}$$

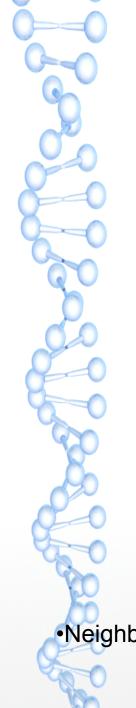
$$d_{ic} + d_{kc} = D_{ik}$$

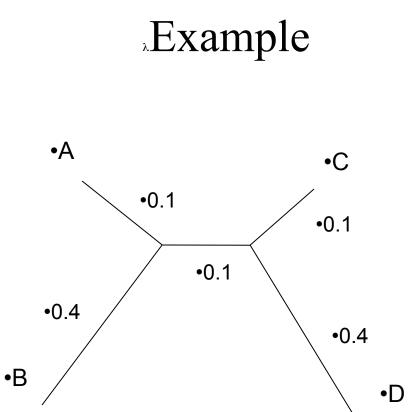
$$d_{jc} + d_{kc} = D_{jk}$$

Four-point condition

Pairwise distances are additive if and only if for every set of four leaves *i,j,k,l*, two of the following three sums are equal and larger than the third:

```
Dij + Dkl
Dik + Djl
Dil + Djk
```





•Neighbor-joining will find the correct tree here

One more thing

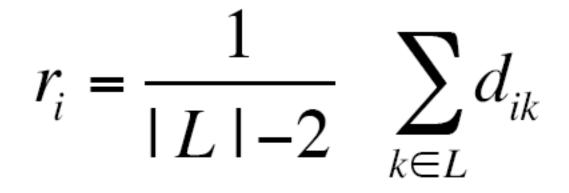
- There are an assortment of formats for trees as there is with DNA sequence data.
- Newick format indicated in the text is one of the more common ones (like FASTA is for sequences)

-Ex: ((A,B),(C,D))

Calculating distances

• Saitou and Nei (1987)

$$D_{ij} = d_{ij} - (r_i + r_j)$$



NJ algorithm

Start with a center star tree and break off pairs

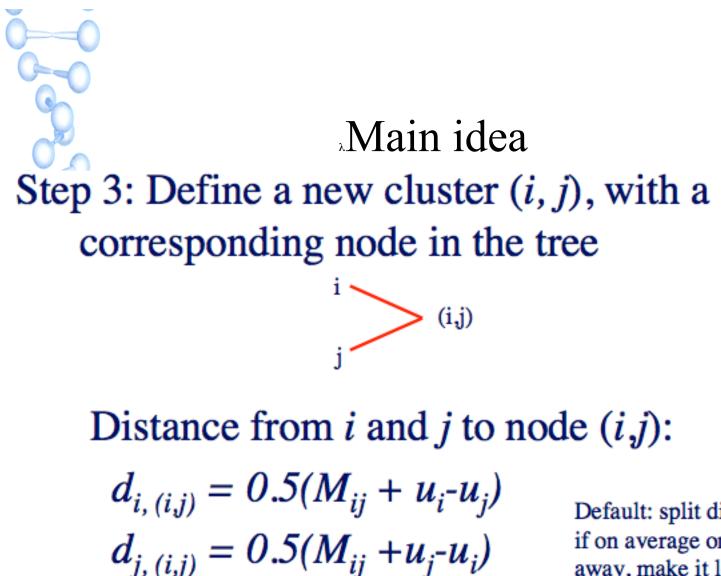
Big caveat:

- Compute (almost) "average" distance to other nodes
- Look for nodes that its distance (Mij) distance from *i* to everything else – distance *j* to everything else is smallest

It turns out the above equates to minimizing branch lengths in the complete tree

Why compute more distances?

- The most crucial piece of NJ is computing a new matrix using the previously mentioned equations.
- This allows us to choose the smallest one greedily, in something called the "4-point condition."
- UPGMA is a simpler form of NJ that is correct when distance between all taxa and the root is the same.
 - This seems true but rarely holds up in observed DNA sequence data



Default: split distance but if on average one is further away, make it longer

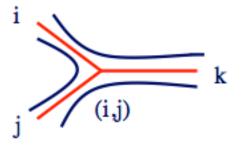




"Finishing up

Step 4: Compute distance between new cluster and all other clusters:

$$M_{(ij)k} = \underline{M}_{ik} + \underline{M}_{jk} - \underline{M}_{ij}$$
2



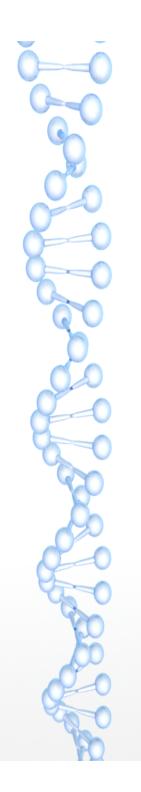
Step 5: Delete *i* and *j* from matrix and replace by (*i*, *j*)

Step 6: Continue until only 2 leaves remain

"Finishing up NJ

Build a NJ tree for the matrix earlier:

	-0	•A	•B	•C	•D	•E
Sa	•A	•0	•5	•3	•8	•10
	•B	•5	•0	•5	•8	•10
Che and a second	•C	•3	•5	•0	•8	•10
	•D	•8	•8	•8	•0	•1
Se	•E	•10	•10	•10	•1	•0



Parsimony

AIntro

Parsimony is a simple and fast approach, making it popular

Two distinct subproblems: Find the history of mutations (easy) Given an alignment, infer tree (hard)

A Parsimony

Intuitively, we want to measure changes along edges of trees.

Similar to Okham's razor: find simplest explanation that works

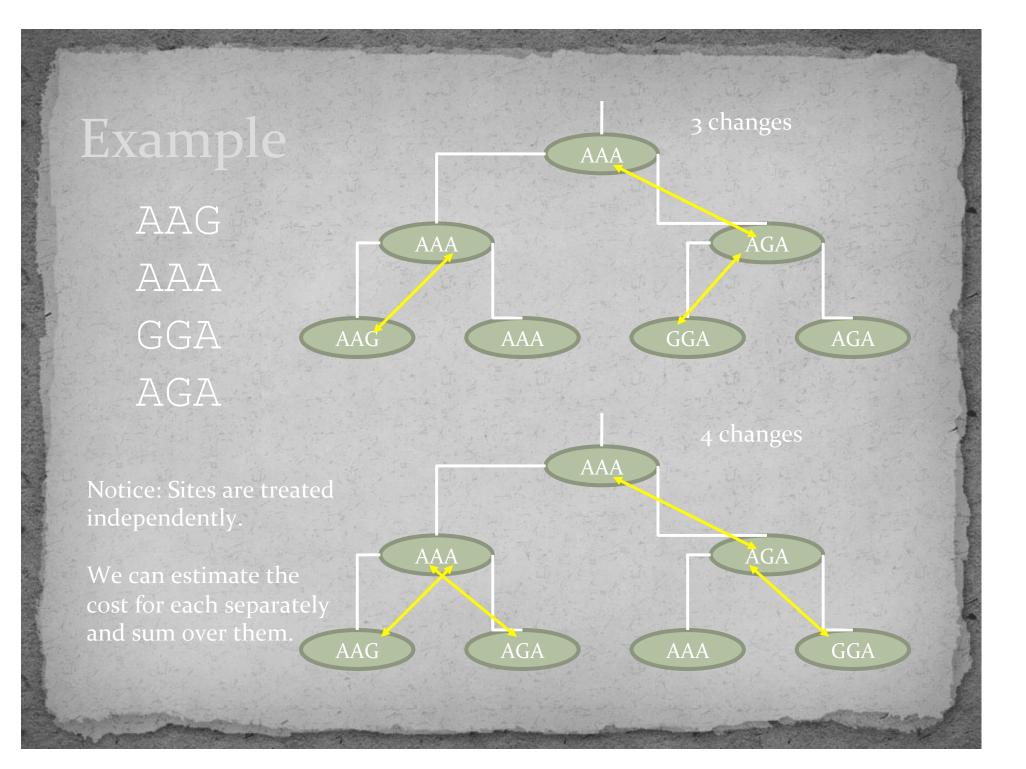
par·si·mo·ny /ˈpärsəˌmōnē/

noun

extreme unwillingness to spend money or use resources.

"a great tradition of public design has been shattered by government parsimony" synonyms: cheapness, miserliness, meanness, parsimoniousness, niggardliness, close-fistedness, closeness, penny-pinching; More





Overview

Input:

Character-based data such as an alignment

Output:

tree that requires minimal number of changes

Goal:

Find right tree topology (how things branch) instead of the actual lengths of edges Hard part is finding optimal topology

^AIn-class example

200

	•1	•2	•3	•4	•5	•6
•Cat	•A	•T	•A	•C	•A	•G
•Dog	•A	•T	•A	•G	•T	•G
•Chim p	•A	•C	•A	•C	•A	•G
•Cow	•T	•C	•G	•G	•T	•A
•Bat	•T	•C	•G	•C	•A	•G

Overview

Input:

Character-based data such as an alignment

Output:

tree that requires minimal number of changes

Goal:

Find right tree topology (how things branch) instead of the actual lengths of edges Hard part is finding optimal topology

"Fitch's algorithm

Published in 1971

Relies on the following assumptions: Any state can convert to any other state Positions in an input are independent Cost is uniform, i.e., A -> T = G -> T

"Fitch's algorithm

We will do two traversals of the tree

Bottom - up (leaves to root) Determine set of possible states for each node

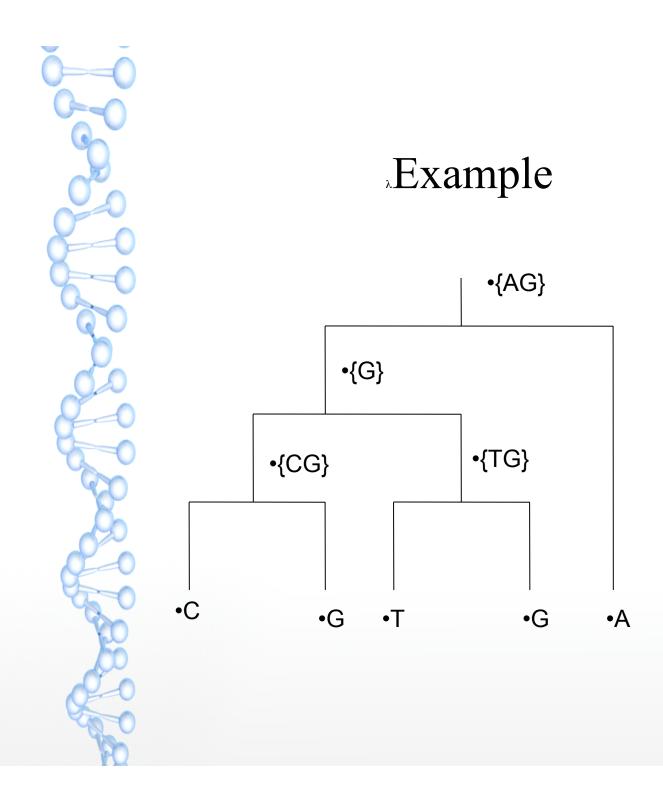
Top - down (root to leaves) Pick the ancestral state for each node from the set of possibilities

»Step 1

Perform a post-order traversal of the tree

Compute:

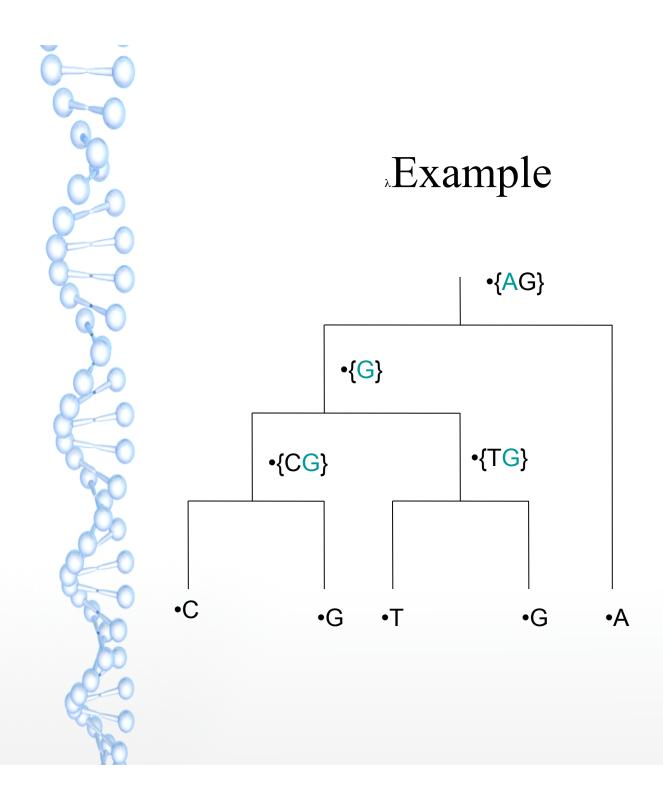
union operations = # of changes



»Step 2

Now do a preorder traversal of tree

Select state *rj* for an internal node *j* with a parent node *i* as follows:



Weighted Parsimony

Sankoff & Cedergren (1983)

Rather than assume all state changes are equally likely, use different costs for different changes

We'll need to propagate costs up during first step of approach, but will not cover this in class

Methods for exploring tree space

Consider any single internal edge in a tree

There are 3 ways the four subtrees can be grouped:

```
AB - CD; AC - BD; AD - BC
```

¹Using this rule to look at trees is called nearest neighbor interchange.

Exact method

There is a branch and bound approach that can be used to calculate the best tree more efficiently.

In short, if a tree you build is worse than a previously discovered tree, you stop

Keep track of all partial trees such that you can reuse information as much as possible

Probabilistic methods

Input:

Given an alignment and a mutation model (e.g., Jukes-Cantor)

Problem:

Compute the likelihood of a tree as a product of the individual likelihoods from the alignment Assumes columns are independent

Conclusions

Many algorithms exist for these problems

Parsimony generally does pretty well for most applications

Likelihood, however, is becoming popular due to increased computational power.

Some future directions for sequencing

Organism sequencing

- Sequence a large fraction of all organisms
- Deduce ancestors
 - Reconstruct ancestral genomes
 - Synthesize ancestral genomes
 - Clone—Jurassic park!

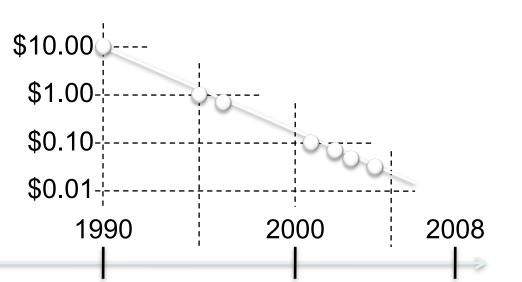
- Study evolution of function
 - Find functional elements within a genome
 - How those evolved in different organisms
 - Find how modules/machines composed of many genes evolved

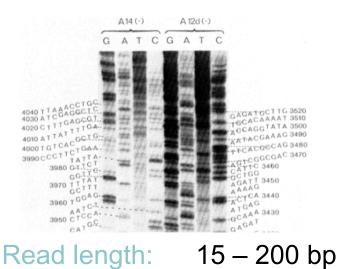
Sequencing technology

Sanger sequencing

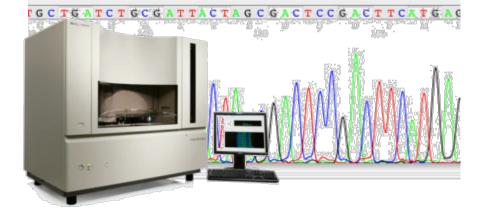
Cost per finished bp:

1975 1980





Throughput:



500 – 1,000 bp

Source bioinformatics.org

"grad-student years"

2 · 10⁶ bp/day



\$985 deCODEme (November 2007)

\$399 Personal Genome Service (November 2007)

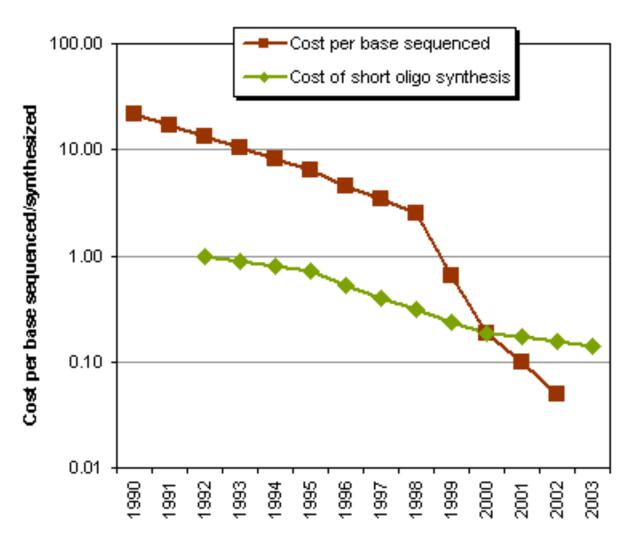


\$2,500 Health Compass service (April 2008)

Genetic Information Nondiscrimination Act (May 2008)



\$~1,000 Whole-genome sequencing (2013)



Sequencing technology

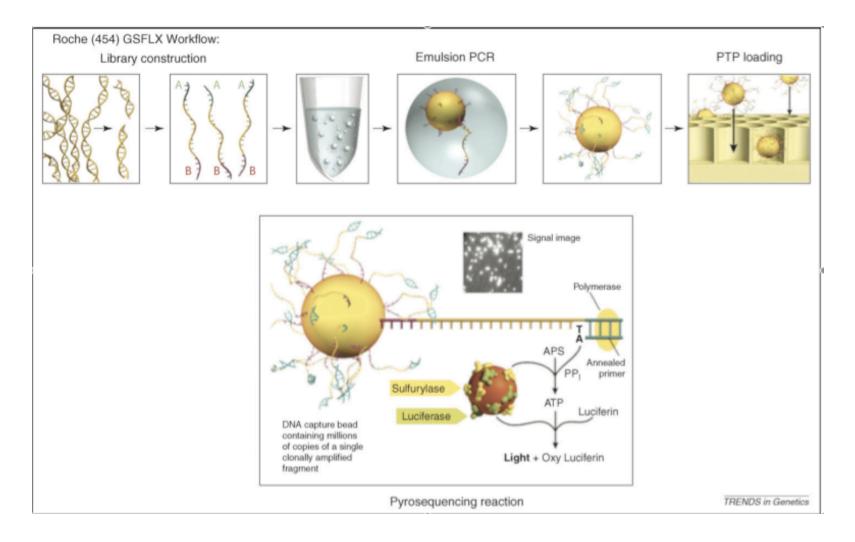
Next-generation sequencing



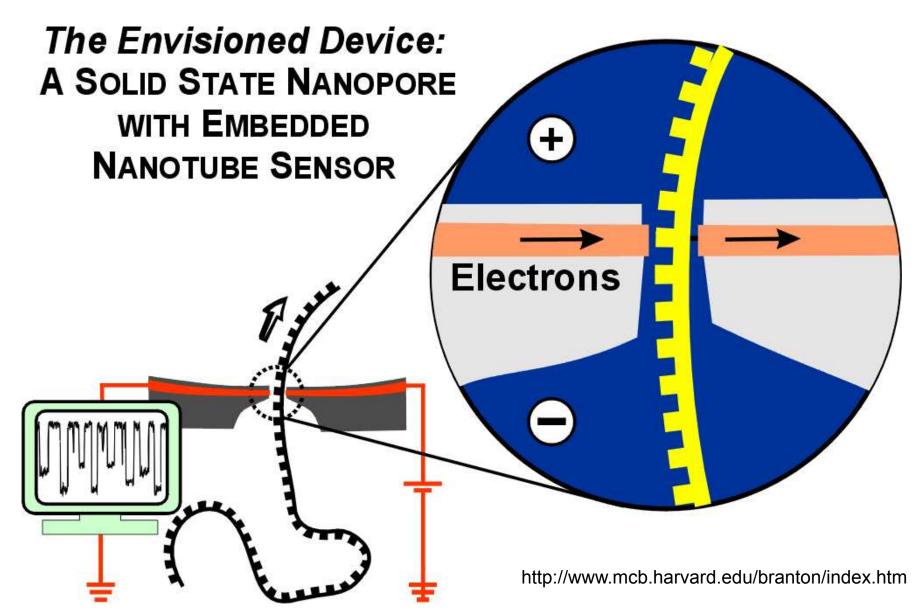
Read length: 450 bp "short reads" Throughput: 600-800 Mb/day Cost: ~ 20,000 bp/\$ De novo: yes

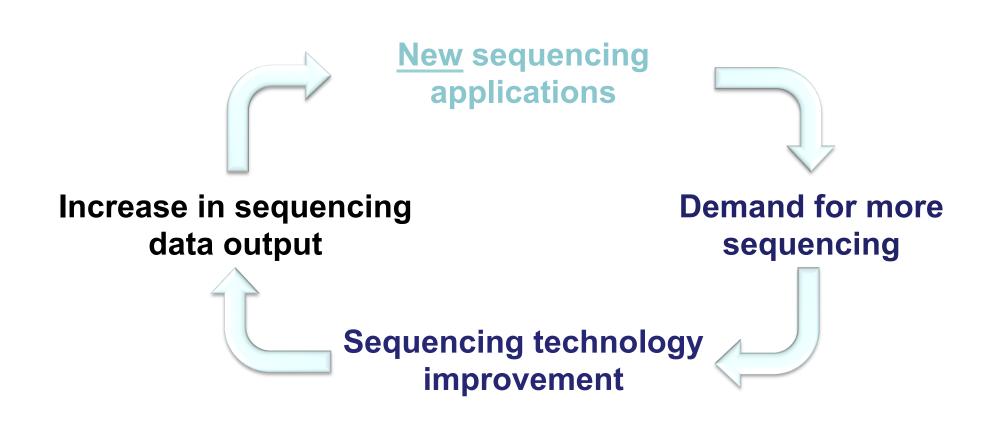
Genome Sequencer / FLX

454 Sequencing



Nanopore Sequencing



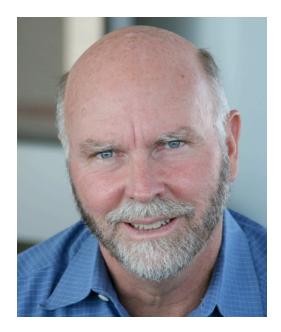


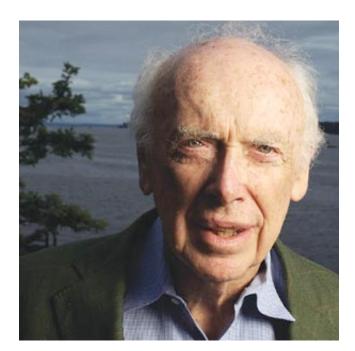
Applications

Whole-genome sequencing Comparative genomics Genome resequencing Structural variation analysis Polymorphism discovery **Metagenomics** Environmental sequencing Gene expression profiling Genotyping Population genetics Migration studies Ancestry inference Relationship inference Genetic screening Drug targeting Forensics

Goal

• Sequencing a human genome





Technology	Read length (bp)	Pairing	bp / \$	de novo
Sanger	1,000	longish	1,000	yes
454	250	shortish	10,000	yes
Solexa/ABI	30	shortish	100,000	maybe

Application	Sanger	454	Solexa/ABI
Bacterial sequencing	yes	yes	maybe
Mammalian sequencing	yes	sort of	probably no
Mammalian <u>re</u> sequencing	Lots of \$\$	Lots of \$\$	yes

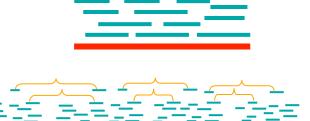
Overlap-Layout-Consensus Assemblers: ARACHNE, PHRAP, CAP, TIGR, CELERA

Overlap: find potentially overlapping reads

Layout: merge reads into contigs and contigs into supercontigs

Consensus: derive the DNA sequence and correct read errors

..ACGATTACAATAGGTT..





Triazzle: A Fun Example

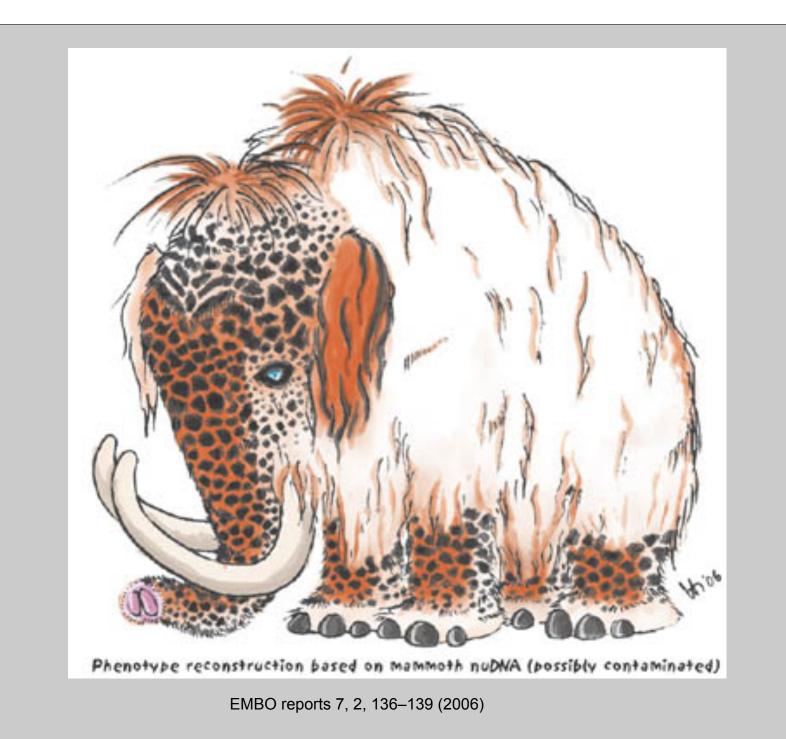
The puzzle looks simple

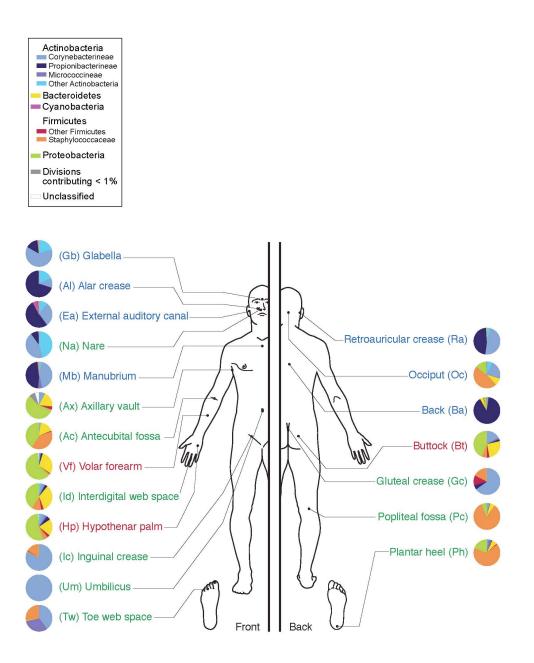
BUT there are repeats!!!

The repeats make it very difficult.

Try it – \$10.99 at www.triazzle.com
 Internet

iPhone version too!

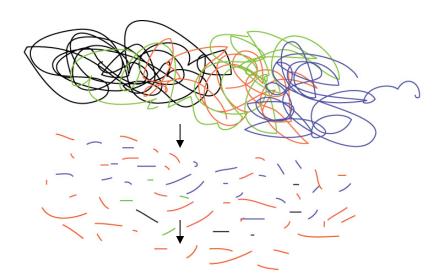






Whole Gene Shotgun Sequencing for Metagenomics

Multiple genomes

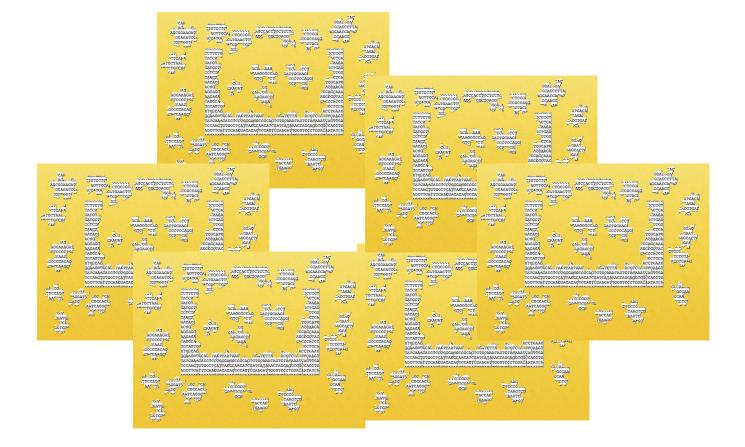


Random genomes fragmentation

...ACACATATACATACAGAGATAGCCCAGATG AGCCCAGATGGCGCTGCTGCTGGGCGCG.....

Genomes assembly using overlaps

...ACACCATACAGAGATAGGGGTGGG GGGTGTGGAGCCCAGATGGCGCTGCTGCTGG......



"Working with thousands of jigsaw puzzles"

What is Metagenomics?

- Metagenomics (Environmental Genomics or Community Genomics) is the study of genomes recovered from environmental samples
- Pro: No need to culture (grow in lab) them
- Con: Heavily uses bioinformatics tools to facilitate insight

Why is Metagenomics Important?

- Some reasons include:
 - Organisms can be studied directly in their environments
 - There are significant advantages for viral metagenomics, because of difficulties cultivating the appropriate host
 - Genomic information has advanced research in a diverse fields such as forensic science

Many projects, many fragments

- Examples:
 - Prokaryote:
 - Sargasso Sea (Venter et al 2004) : **1.6 billion** base pairs generated estimated to come from **1800** genomic species
 - Viral:
 - Marine water (Breitbart et al 2002) Mission Bay and Scripps Pier. 873 sequences for the Mission Bay and 1061 for Scripps Pier with respectively more than 65% and 73% of unknown

Many projects, many fragments

- Three years after the Marine Water project, most of sequences are still unique. Despite the fact that GenBank has more than doubled in size.
- All of the Metagenome projects have generated enormous amounts of data that still cannot be assembled or annotated.