Sequence alignment (continued)

#### Various types



### Global Alignment (review)

- Input: Two strings, labeled s and t
  - |*s*| is *n*
  - |*t*| is *m*
- Output: Two strings  $s_A$  and  $t_A$  such that:
  - $-s_A$  and  $t_A$  are of equal length L
  - Characters must be in same order, with "-" spacers as needed
  - If  $s_A[i] =$  '-' , then  $t_A[i] \neq$  '-'
  - If  $t_A[i] =$  '-' , then  $s_A[i] \neq$  '-'

#### Visual alignments

#### Score: A measure of alignment quality

|            | С         | Α      | Т    | -    | Т      | С    | Α   | _    | С     |        |      |
|------------|-----------|--------|------|------|--------|------|-----|------|-------|--------|------|
|            | С         | —      | Т    | С    | G      | С    | Α   | G    | С     |        |      |
| -          | 10        | <br>-5 | 10   | -5   | <br>-2 | 10   | 10  |      | 10    |        |      |
| Tota       | al =      | 3      | 3    |      |        |      |     |      |       |        |      |
| Scor<br>et | red<br>c. | as     | E (C | :,C) | E(     | (A,- | -), | Е (І | Γ,Τ), | , E(-, | ,C), |

# Types of scores

- Reasons:
  - Sequencing error(s)
  - Evolutionary change
- Three parameters
  - Gap (indel?)
  - Mismatch (misread base?)
  - Match (no change)
- We will only consider constant gap penalties for now.

#### **Optimal alignment**

- We want to find an alignment that optimizes an evaluation function  $E(s_A, t_A)$
- The "brute force" method is a hard problem (NP-hard) requiring exponential time.
- Dynamic programming, in cases we will discuss, can be used to solve this problem efficiently.

#### Requirements

- We will need four things to compute a global alignment:
  - 1. Substitution matrix (parameters)
  - 2. Recurrence relation
  - 3. Filling up a table
  - 4. Traceback

We can solve this based on looking at three smaller problems

$$T[i,j] = \max \begin{cases} T[i-1,j-1] + score(s[i],t[j]) \\ T[i-1,j] + g \\ T[i,j-1] + g \end{cases}$$





+10 for match, -2 for mismatch, -5 for space (rowwise)



| λ | 0   | -5  | -10  | -15 | -20    | -25 | -30 | -35 | -40 |
|---|-----|-----|------|-----|--------|-----|-----|-----|-----|
| C | -5  | 10  | 5    | 0   | -5     | -10 | -15 | -20 | -25 |
| A | -10 | 5   | 8    | 3   | -2     | -7  | 0   | -5  | -10 |
| Т | -15 | 0   | 15   | 10  | 5<br>* | 0   | -5  | -2  | -7  |
| Т | -20 | -5  | 10 * | 13  | 8      | 3   | -2  | -7  | -4  |
| C | -25 | -10 | 5    | 20  | 15     | 18  | 13  | 8   | 3   |
| A | -30 | -15 | 0    | 15  | 18     | 13  | 28  | 23  | 18  |
| C | -35 | -20 | -5   | 10  | 13     | 28  | 23  | 26  | 33  |

Traceback yields both optimal alignments in this example

## End-gap free alignment

- We often don't want to penalize gaps at the start or end of the alignment, especially when comparing short and long sequences
- Same as global alignment, except:
  - Initialize with zeros (free gaps at start)
  - Locate max in the last row/column (free gaps at end)
- Also called *semiglobal alignment*

#### Example paths



#### Example traceback paths



-AGCC-TAG-CCAA

AGCTAAG-CT-A-A

AGCCTA-GCCAA

AGCT-AAGCTAA

AGCCTA-GCCAA AGC-TAAGCTAA

--AGCCTAGCCAA \_\_\_\_\_ AGCTAAGCTAA-----

#### Practice

T[i, 0] = 0 T[0, j] = 0

$$T[i,j] = \max \begin{cases} T[i-1,j-1] + score(s[i],t[j]) \\ T[i-1,j] + g \\ T[i,j-1] + g \end{cases}$$

Match = 5 Mismatch = -2 Gap = -3

S = CATTAT = ACATTAG

Find the score of the best semiglobal alignment



+10 for match, -2 for mismatch, -5 for space (rowwise)

#### Local Alignment

T[i, j] = Score of optimally aligning a suffix of *s* with a suffix of *t*.

$$T[i,j] = \max \begin{cases} T[i-1,j-1] + score(s[i],t[j]) \\ T[i-1,j] + g \\ T[i,j-1] + g \\ 0 \end{cases}$$

Initialize top row and leftmost column to zero.

#### Practice

T[i, 0] = 0 T[0, j] = 0  $T[i, j] = \max \begin{cases} T[i-1, j-1] + score(s[i], t[j]) \\ T[i-1, j] + g \\ T[i, j-1] + g \\ 0 \end{cases}$ Matc

S = CTACTT = ATACG Match = 5 Mismatch = -2 Gap = -3

Find the score of the best local alignment



+1 for a match, -1 for a mismatch, -5 for a gap

#### Reducing space requirements

- O (*mn*) tables are often the limiting factor in computing large alignments
- There is a linear space technique that only doubles the time required [Hirschberg77]

# Alignment and new architectures

 http://en.wikipedia.org/wiki/Smith-Waterman\_algorithm